NETWORK-BASED COMPUTING LABORATORY

Designing Multi-Leader based Allgather Algorithms for Multi-core Clusters

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Outline

- Introduction and Background
- Motivation
- Related Work
- Multi-Leader based Algorithms
- Experimental evaluation
- Conclusions and Future Work

Introduction and Background

- MPI is the de-facto programming model for HPC
- Multi-core clusters are becoming increasingly common
- Modern interconnects like InfiniBand offer high-bandwidth and low-latency
- The collective communication primitives consume a significant amount of time
- Necessary to have multi-core aware collective designs

Allgather Communication

- Each process broadcasts a vector data to every other process in the group
- Commonly used algorithms :
  - Recursive Doubling (RD) Algorithm for small messages

    $$tcomm = ts * log(p) + tw * (p - 1) * m$$

  - Ring Algorithm for large messages

    $$tcomm = (ts + tw * m) * (p - 1)$$

tcomm - *Total Communication cost*

ts - *Communication start-up cost*

tw - *Cost of sending a byte of data*

p - *Number of processes*

m - *Message Size.*

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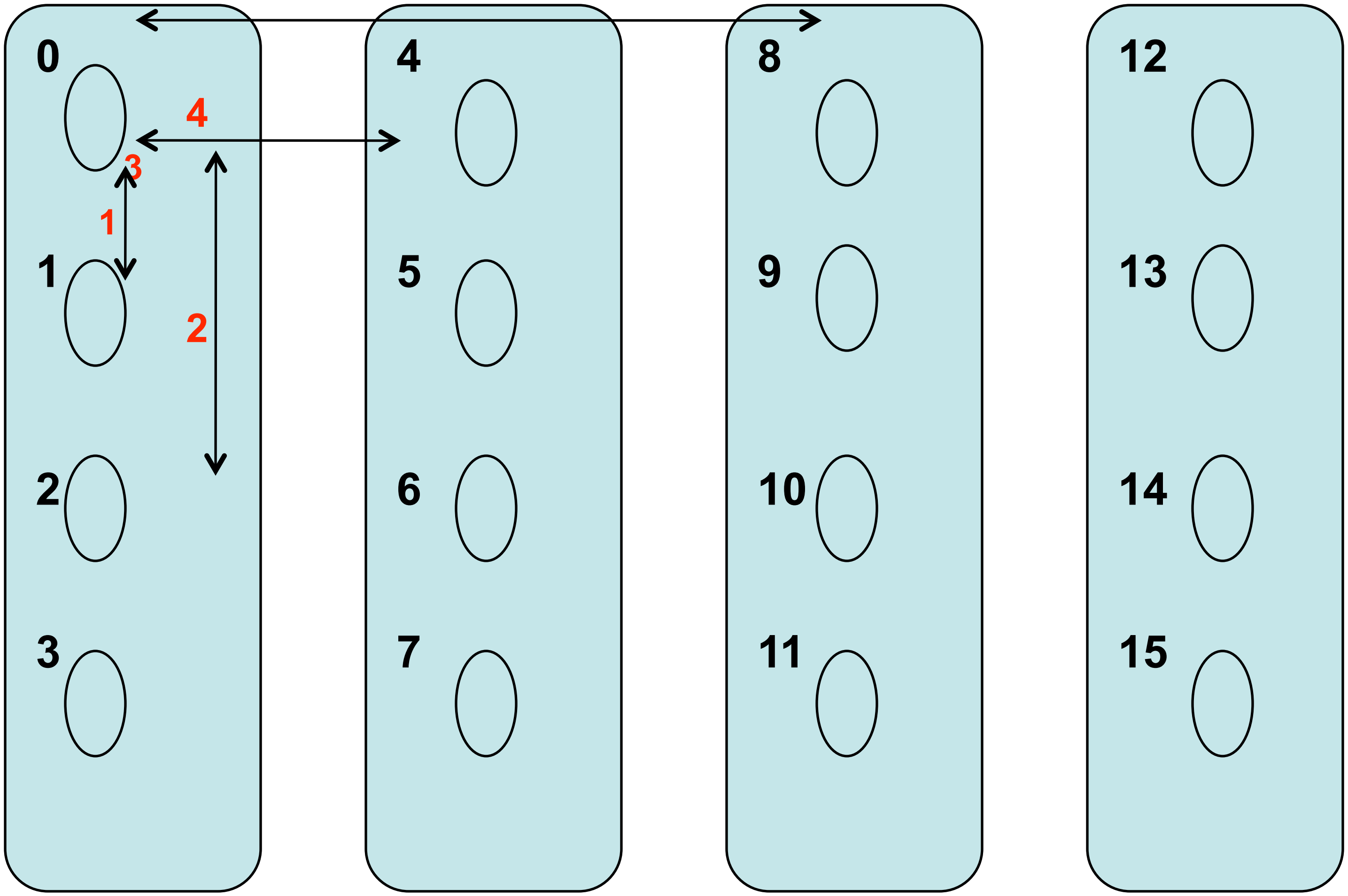
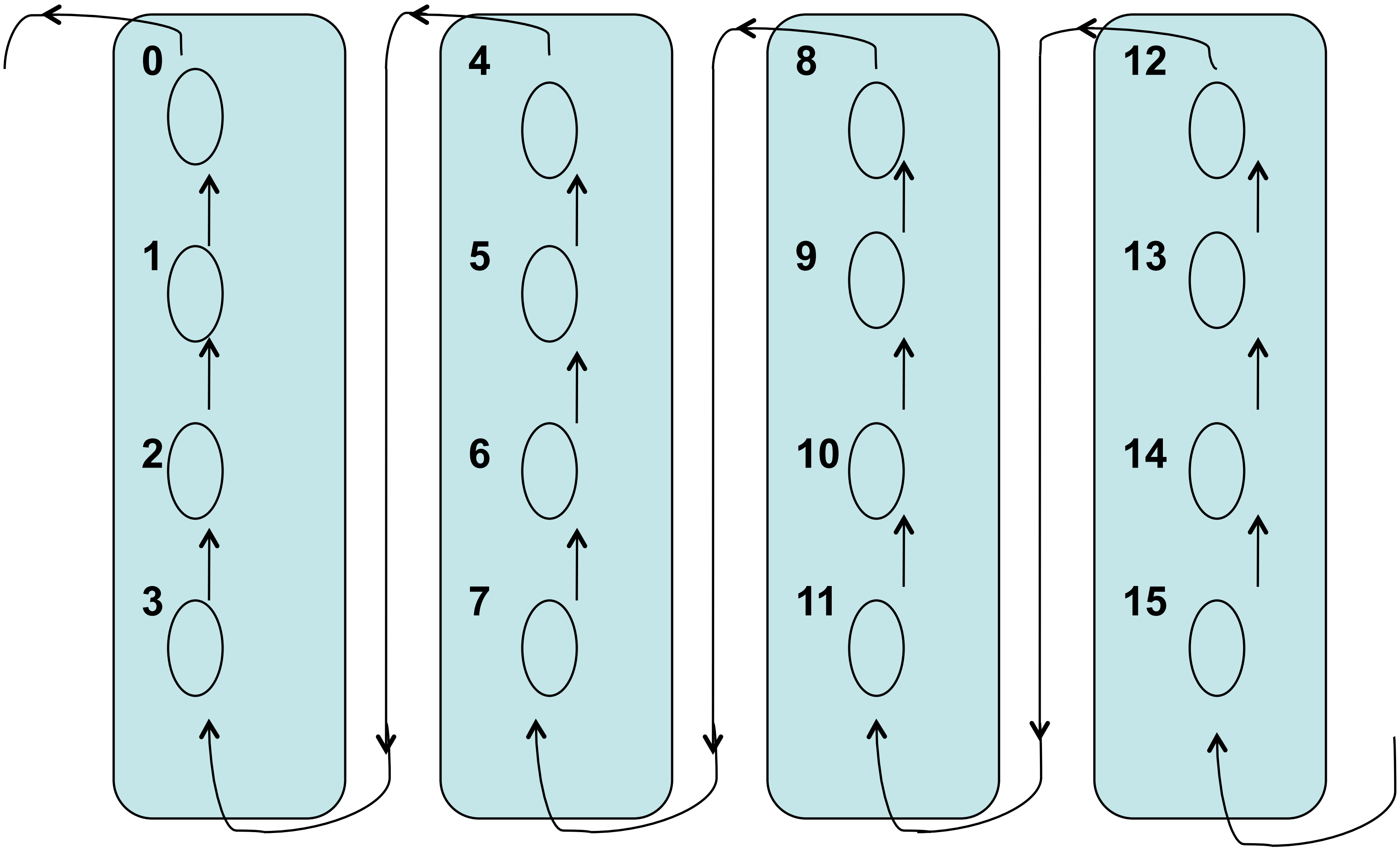
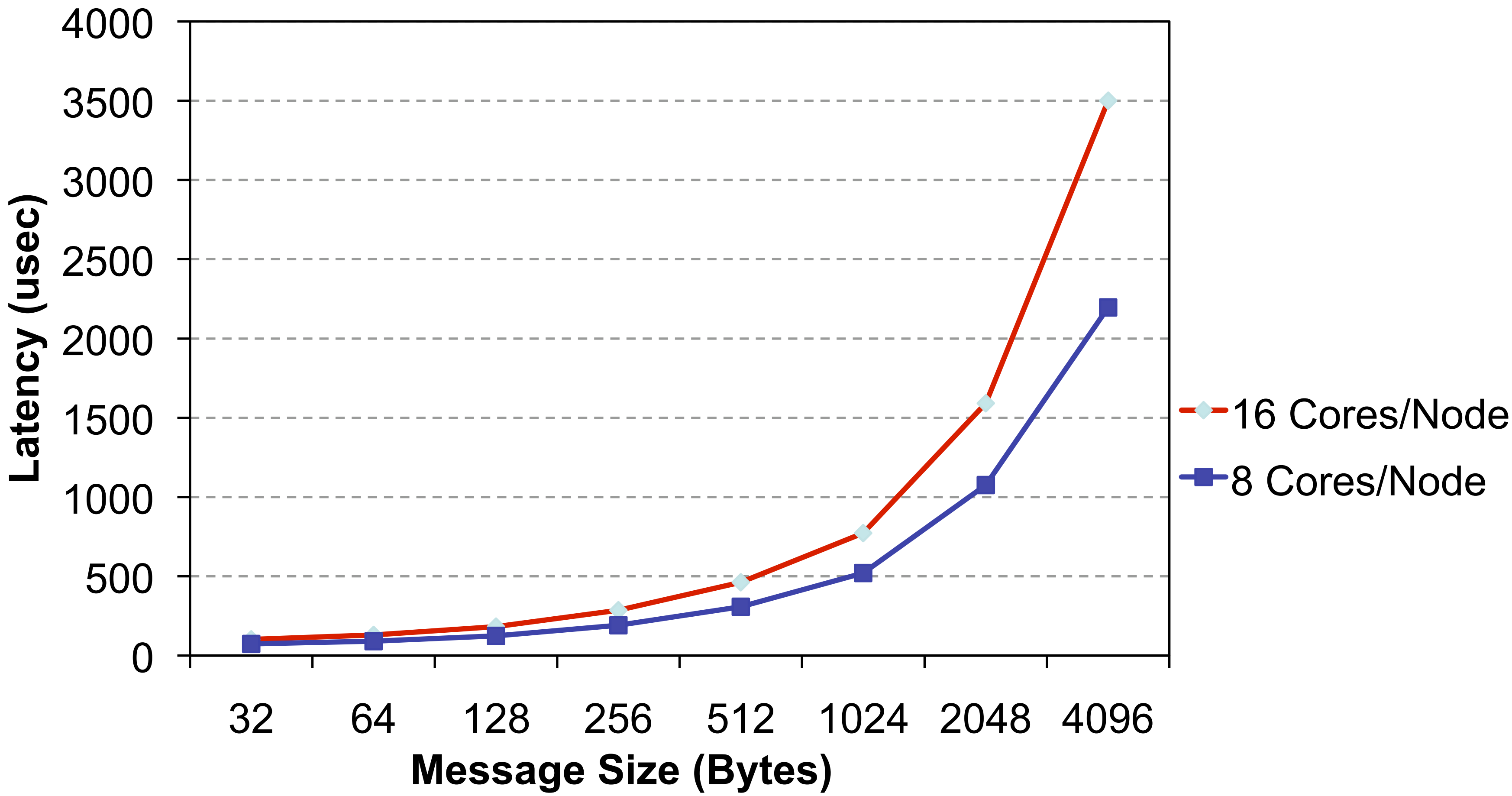
Recursive Doubling (RD) Algorithm on Multi-cores



Ring Algorithm on Multi-cores

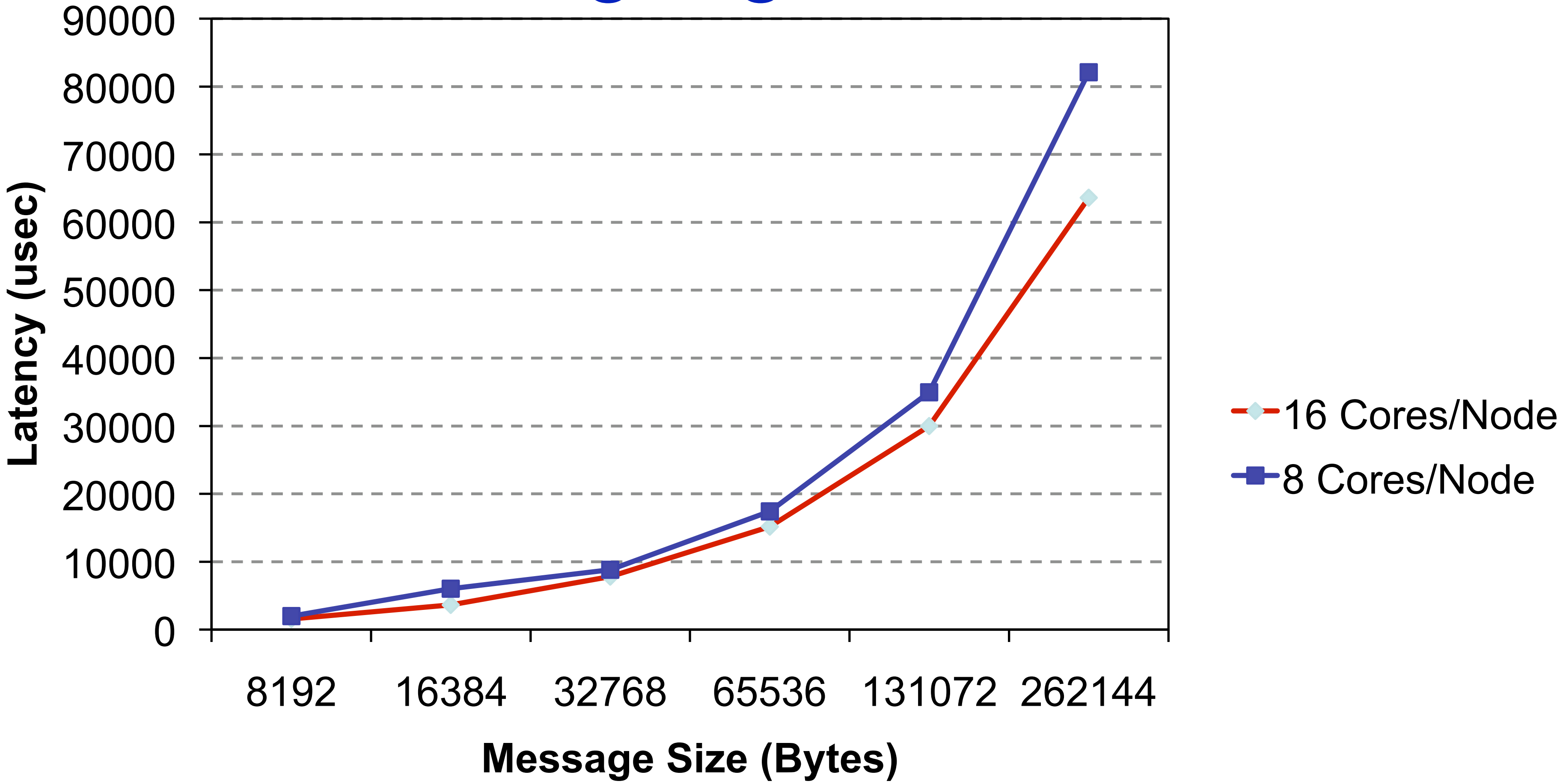


Scaling on Multi-cores : Recursive Doubling Algorithm



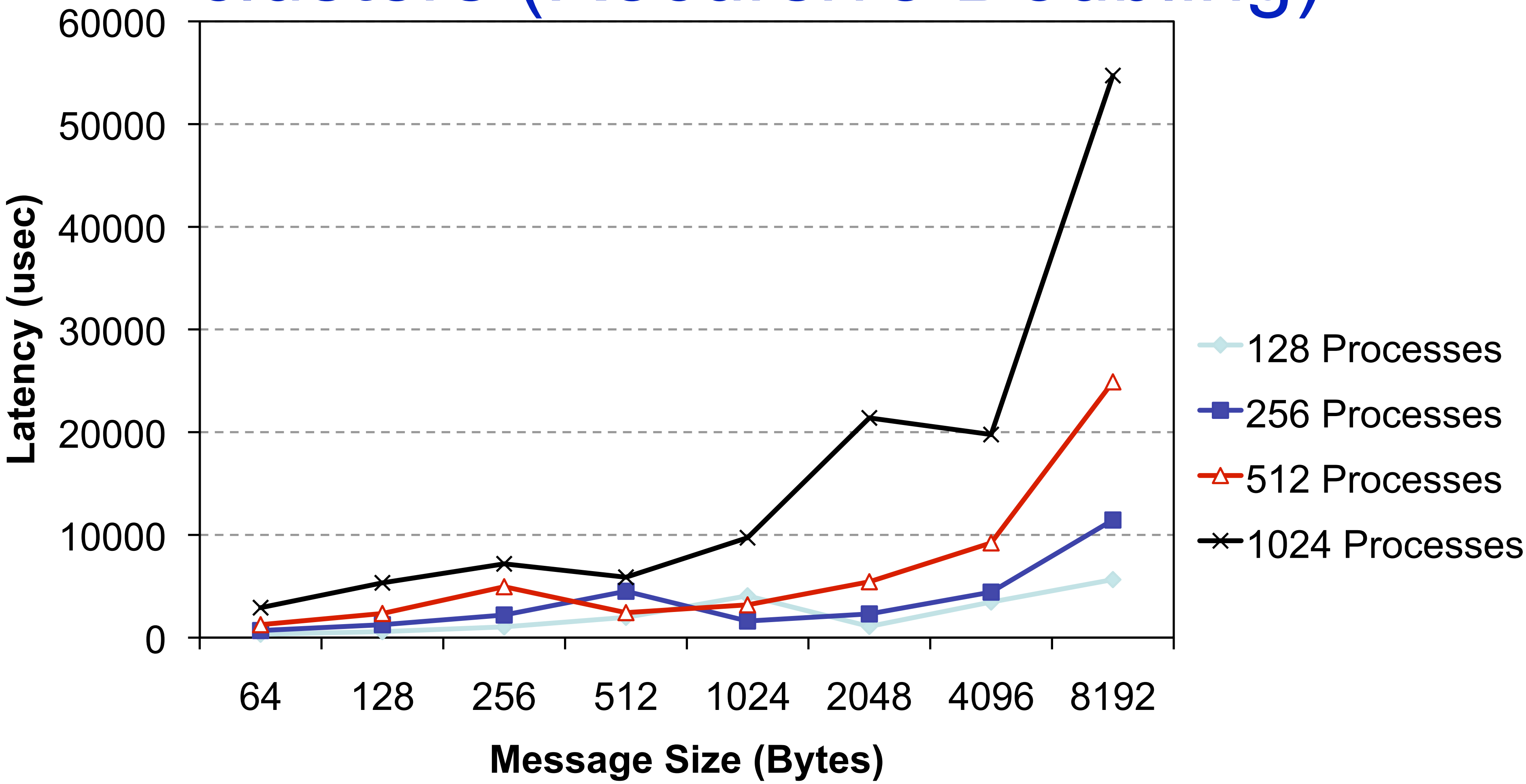
Recursive Doubling (RD) scales poorly with increasing core counts

Scaling on Multi-cores : Ring Algorithm



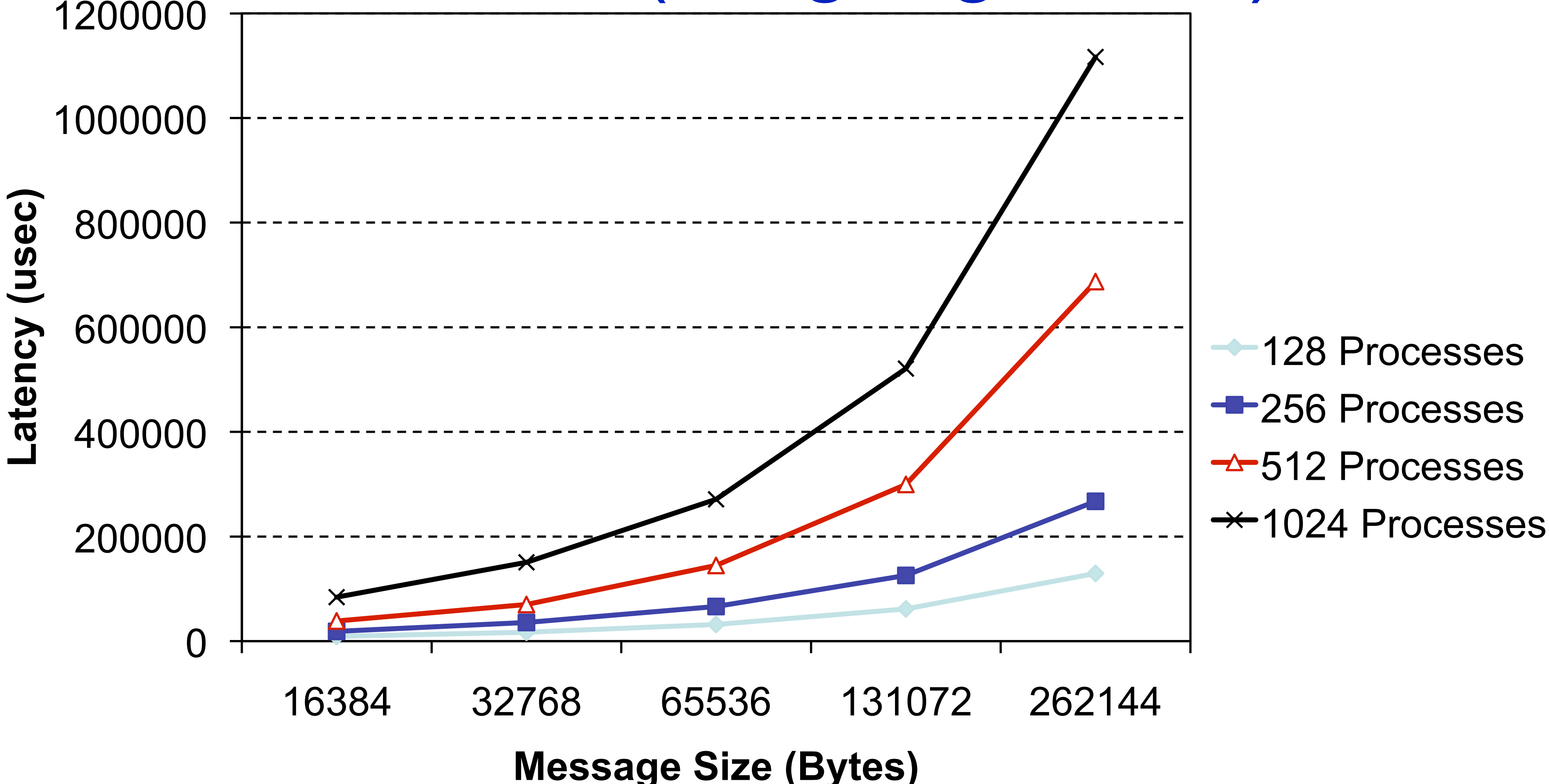
Ring Algorithm scales as expected with increasing core counts

Scaling on Large Scale Multi-core clusters (Recursive Doubling)



Recursive Doubling (RD) scales poorly for large system size

Scaling on Large Scale Multi-core clusters (Ring Algorithm)



Ring Algorithm scales as expected for large system sizes

Problem Statement

- Is it possible to design an algorithm to :
- be Multi-core and NUMA aware to achieve better performance and scalability as core-counts and system sizes increase?
- fully exploit the differential memory access costs in NUMA based Multi-core systems?

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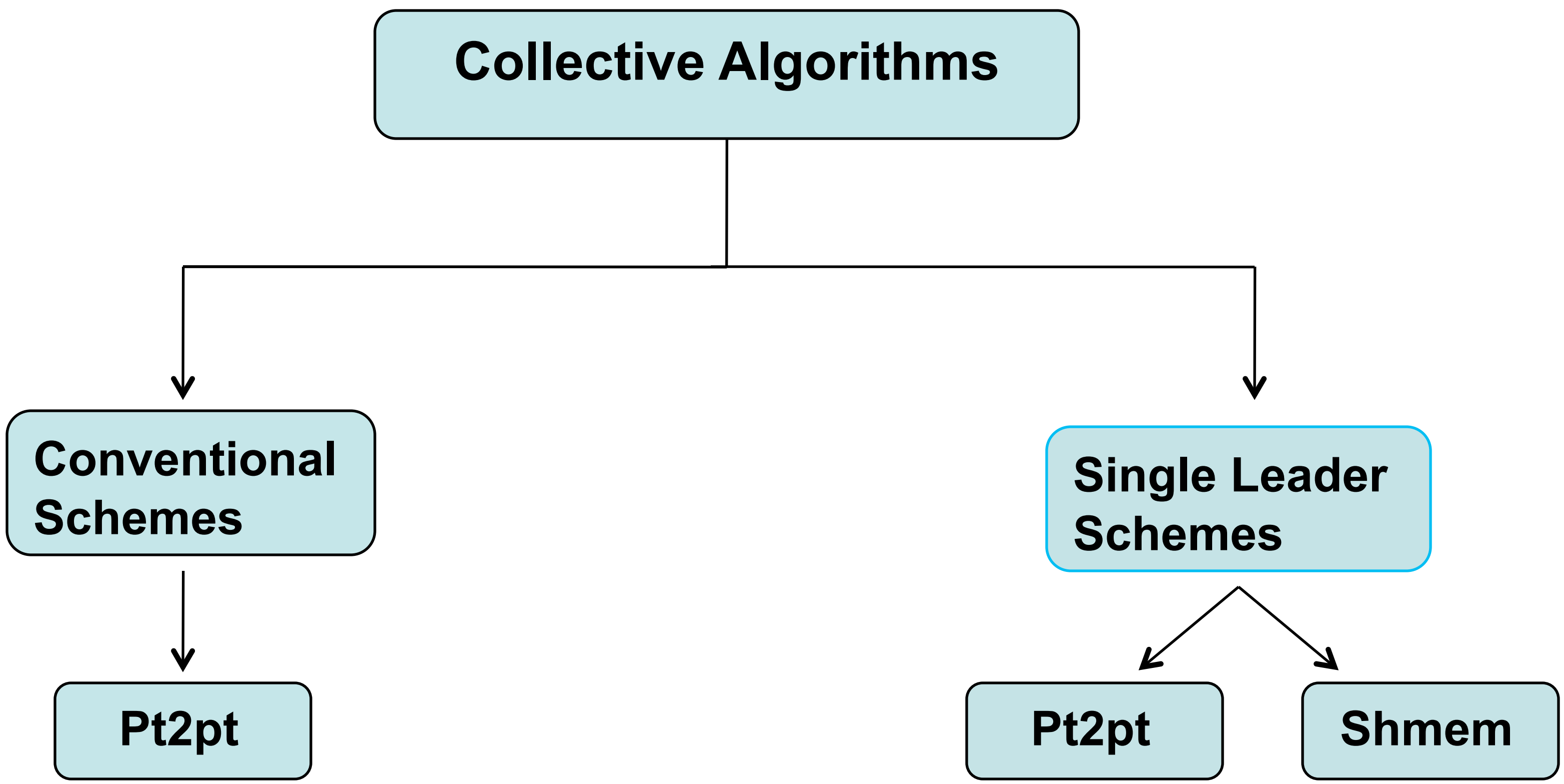
Collective Design Framework



Existing Multi-core aware Algorithms

- Single Leader approaches :

  Aggregation – Distribution .

  Step 1 : Data aggregation at the leader on each node

  Step 2 : Inter leader exchanges

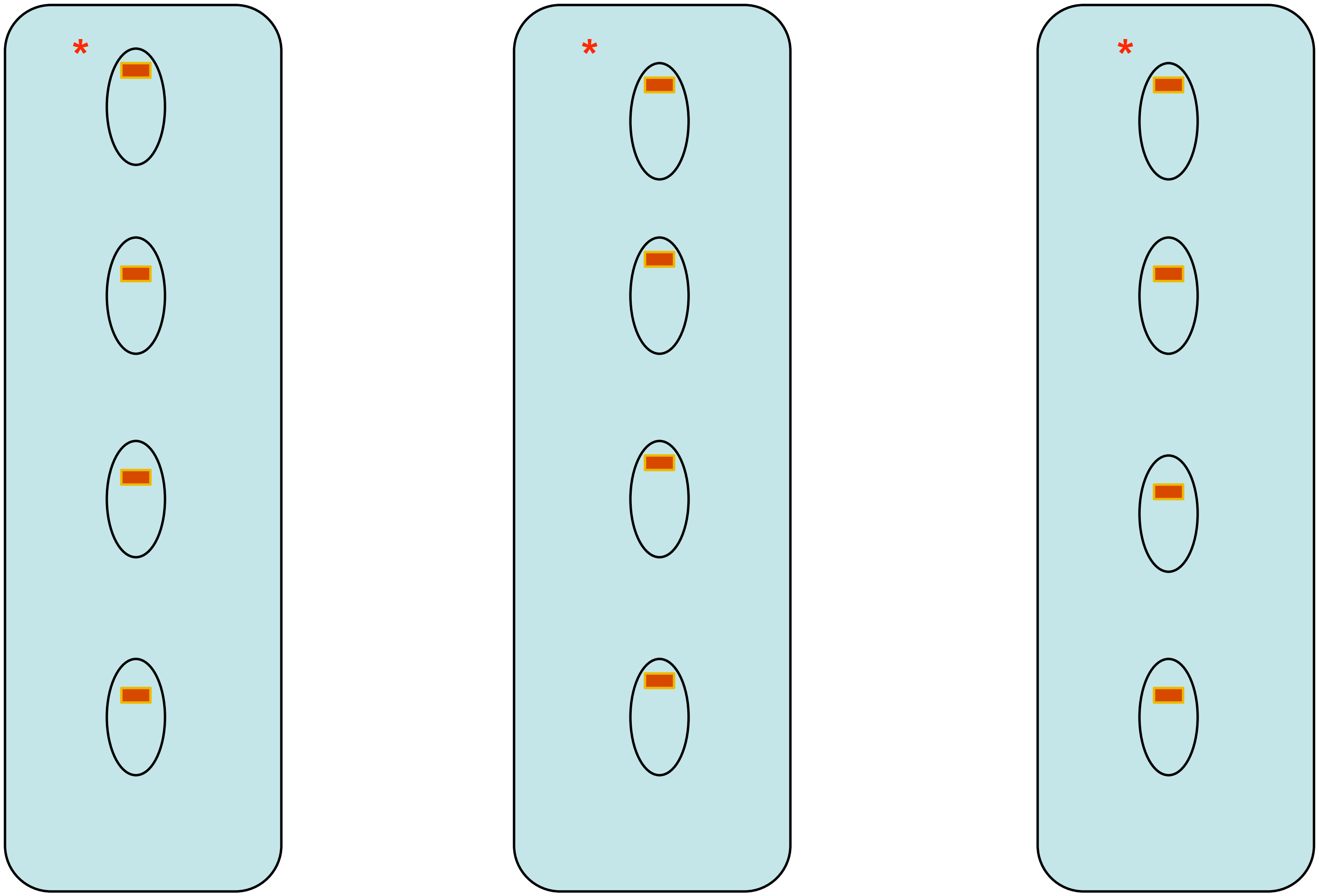
  Step 3 : Data distribution within each node

  Steps 1 and 3 are intra-node operations.

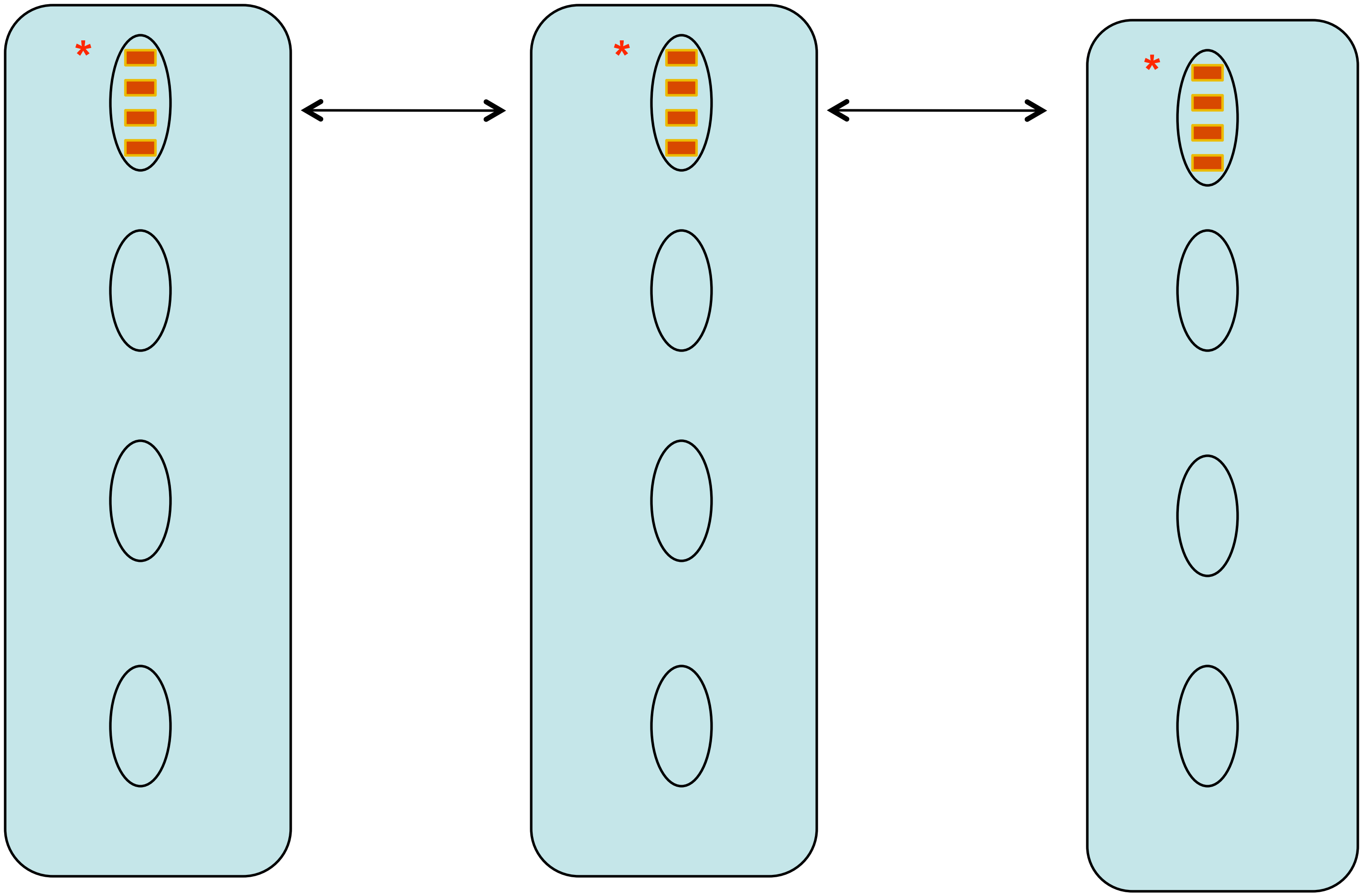
  → Point-to-point MPI calls

  → Shared memory buffer visible to all the processes within a node

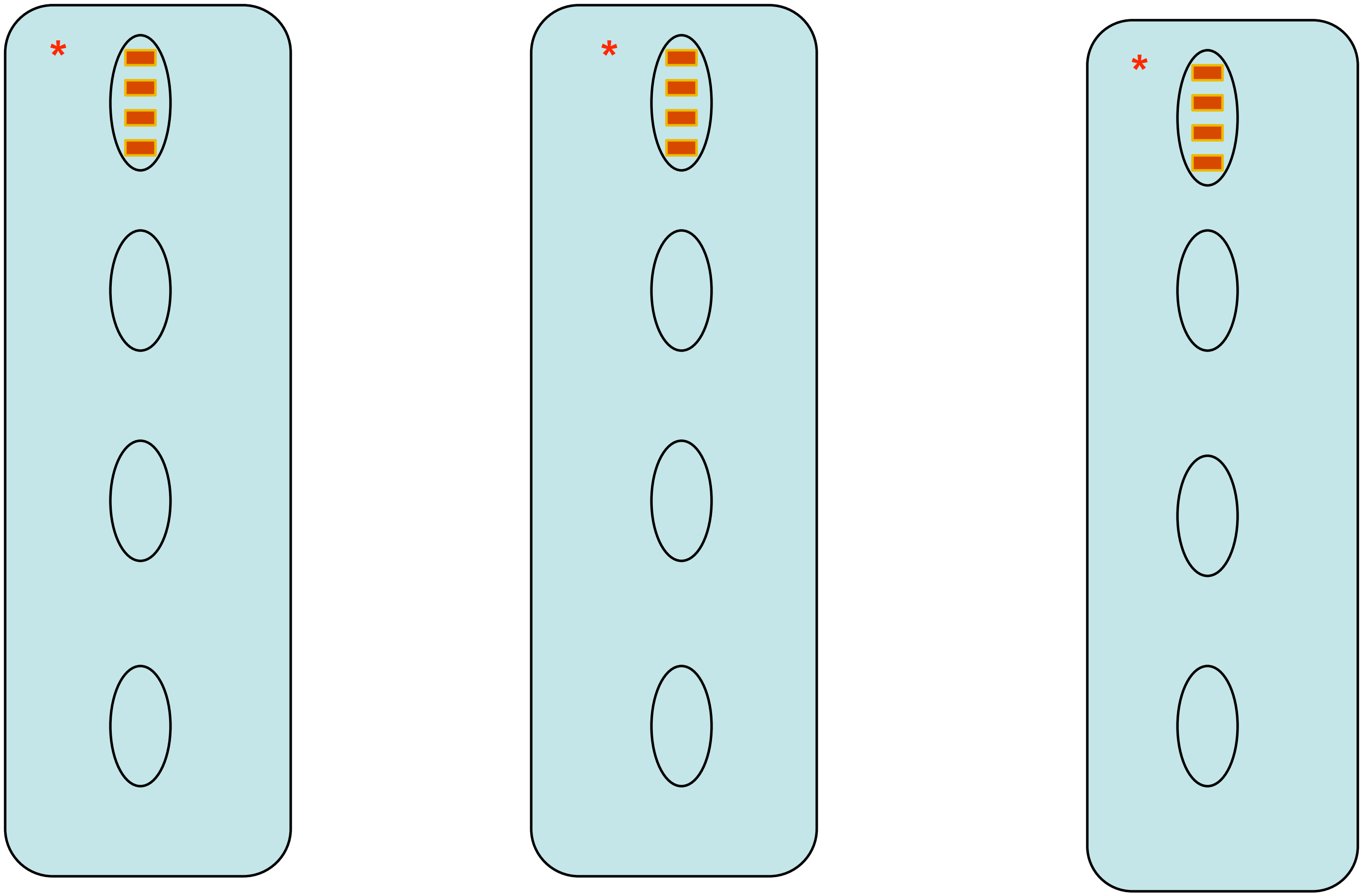
Single Leader Algorithms : Step1 intra-node (pt2pt)



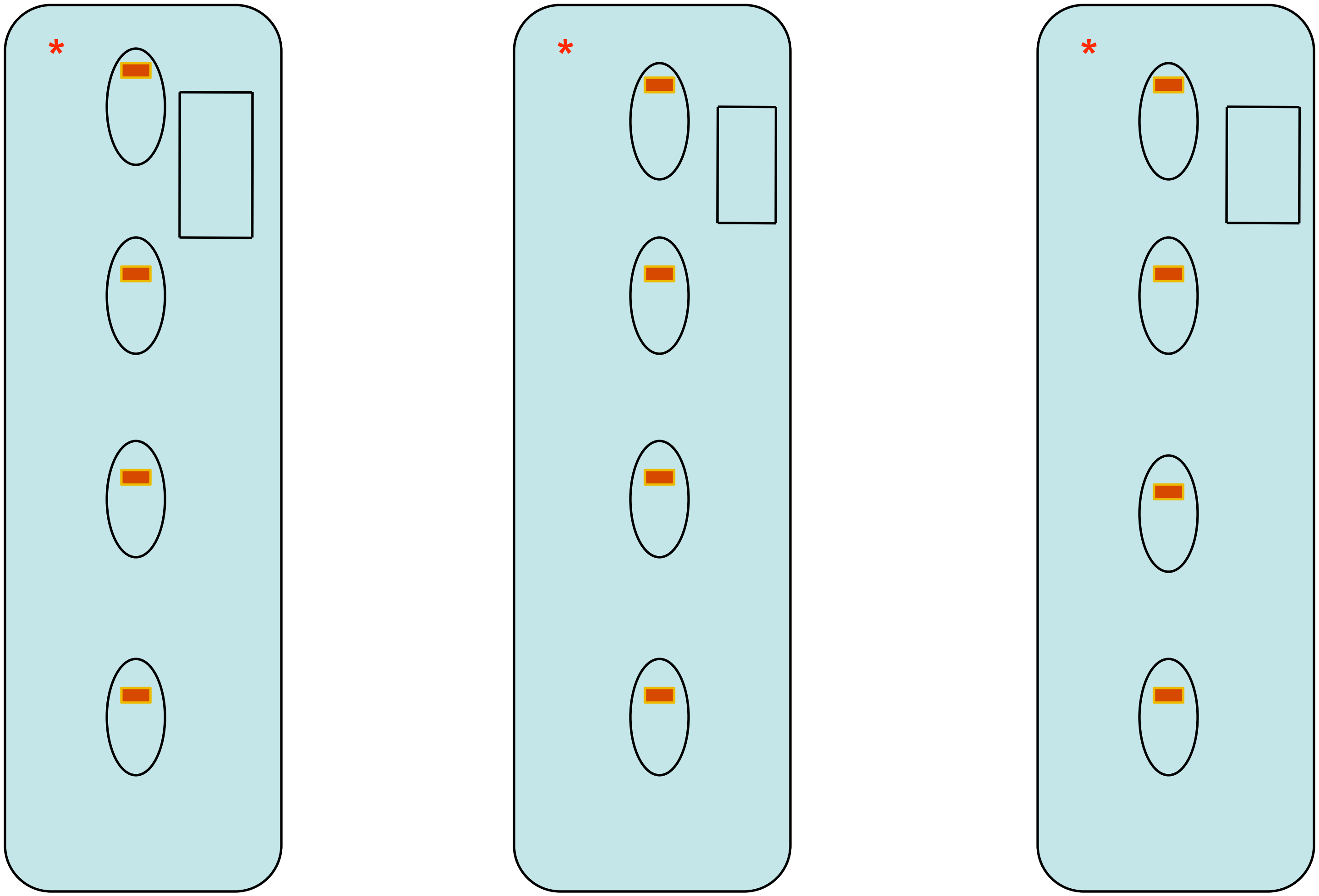
Single Leader Algorithms : Step2 inter-node (pt2pt)



Single Leader Algorithms : Step 3 intra-node (pt2pt)



Single Leader Algorithms : Step1 intra-node (shmem)



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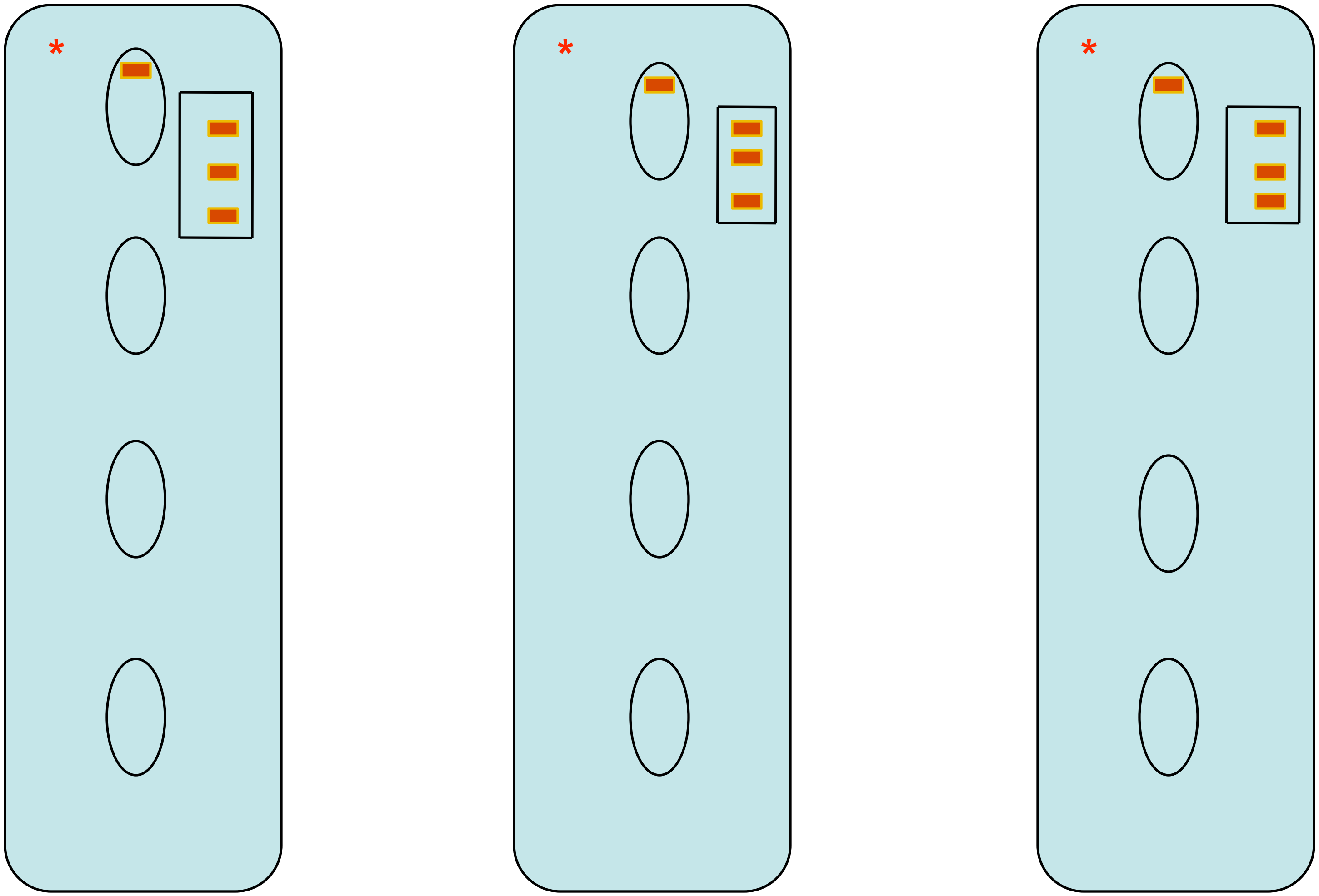
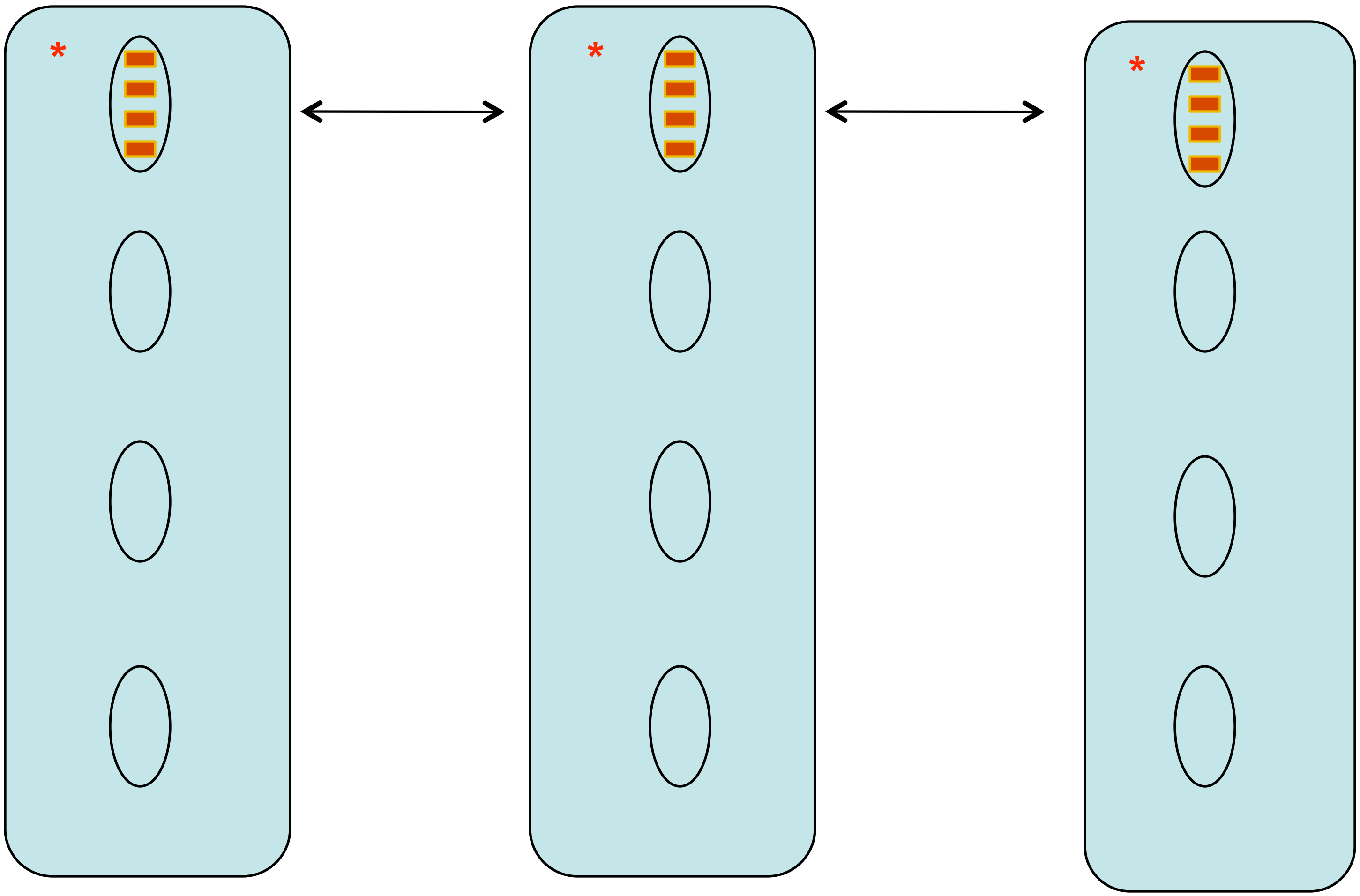


Single Leader Algorithms : Step2 inter-node (pt2pt)



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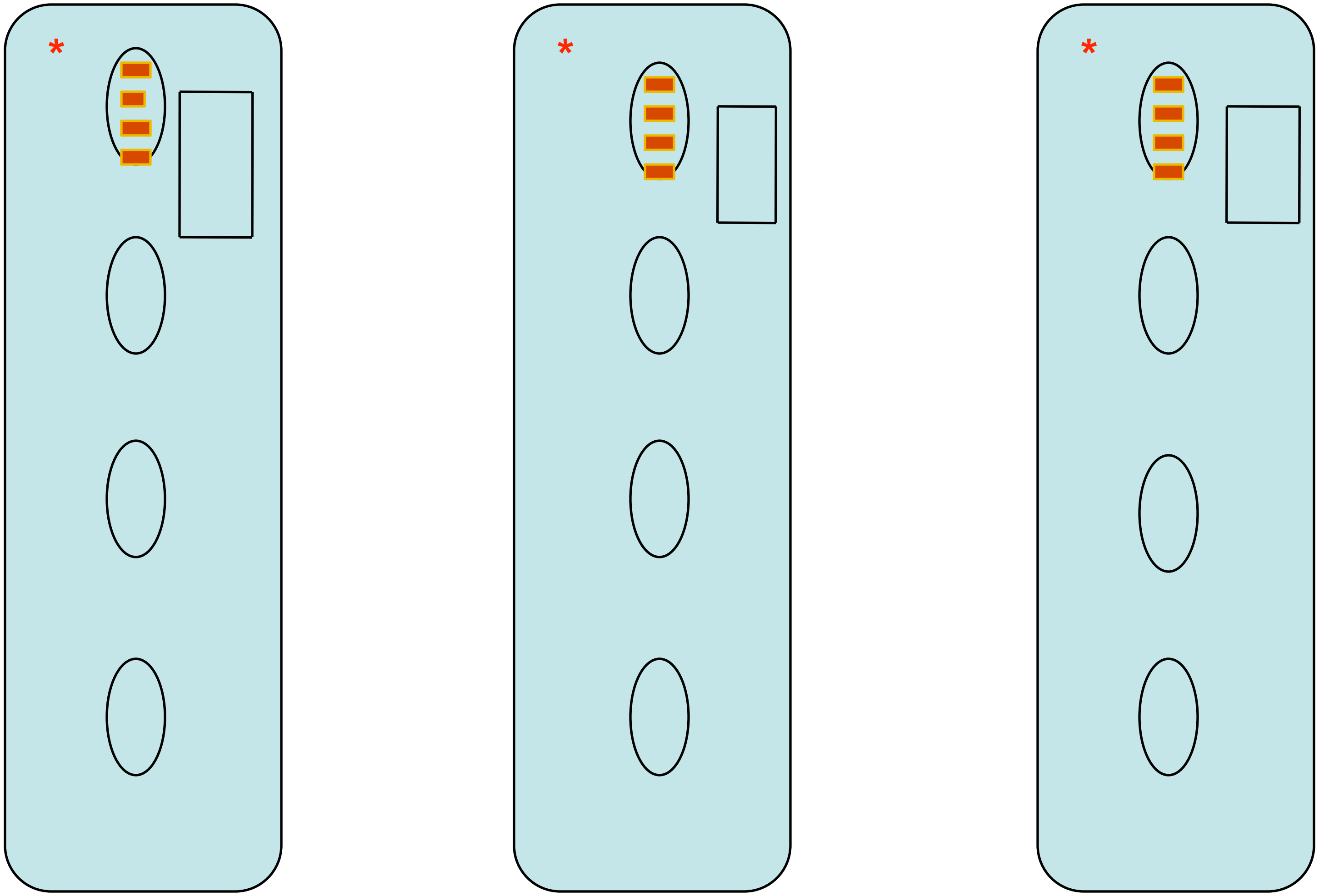
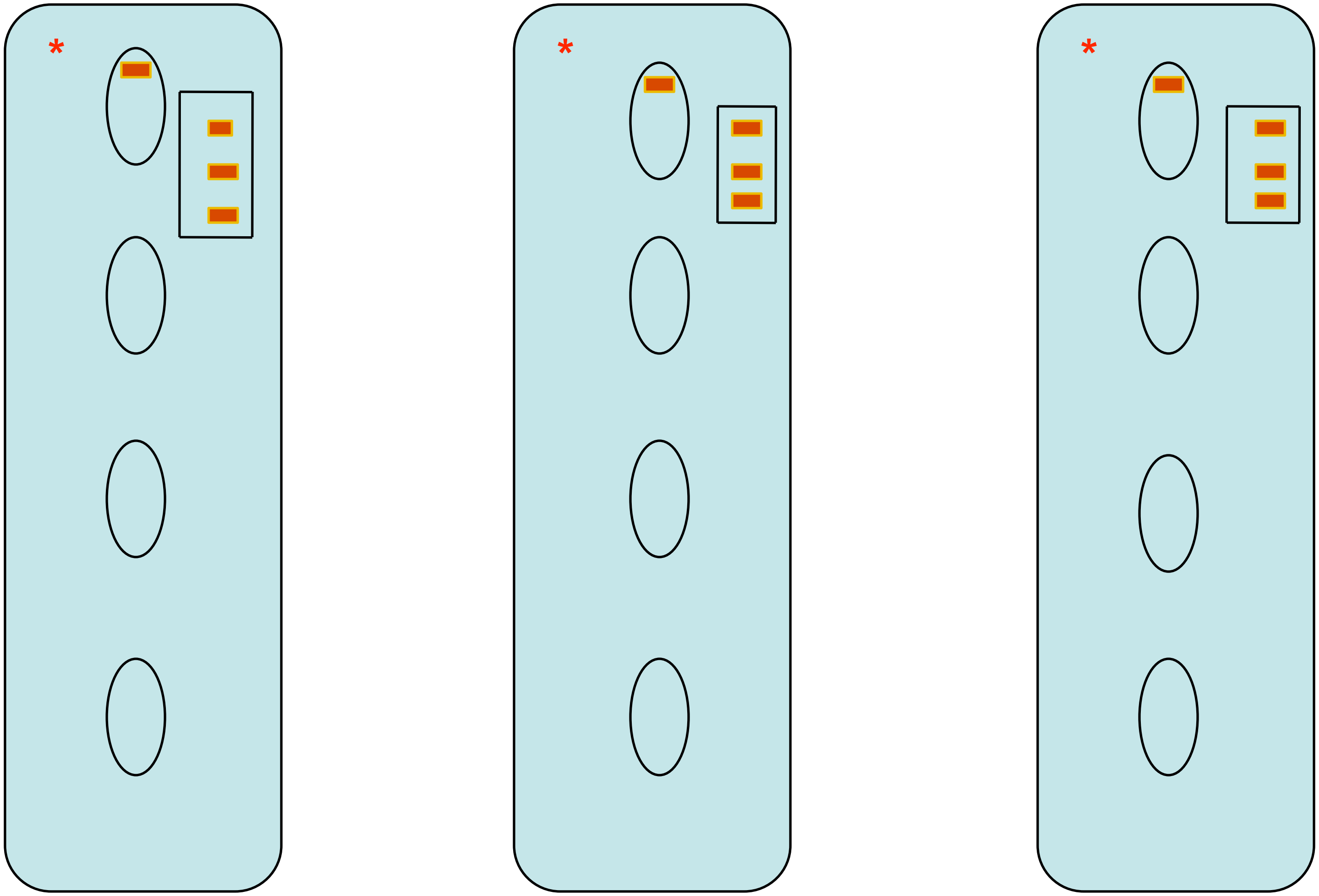
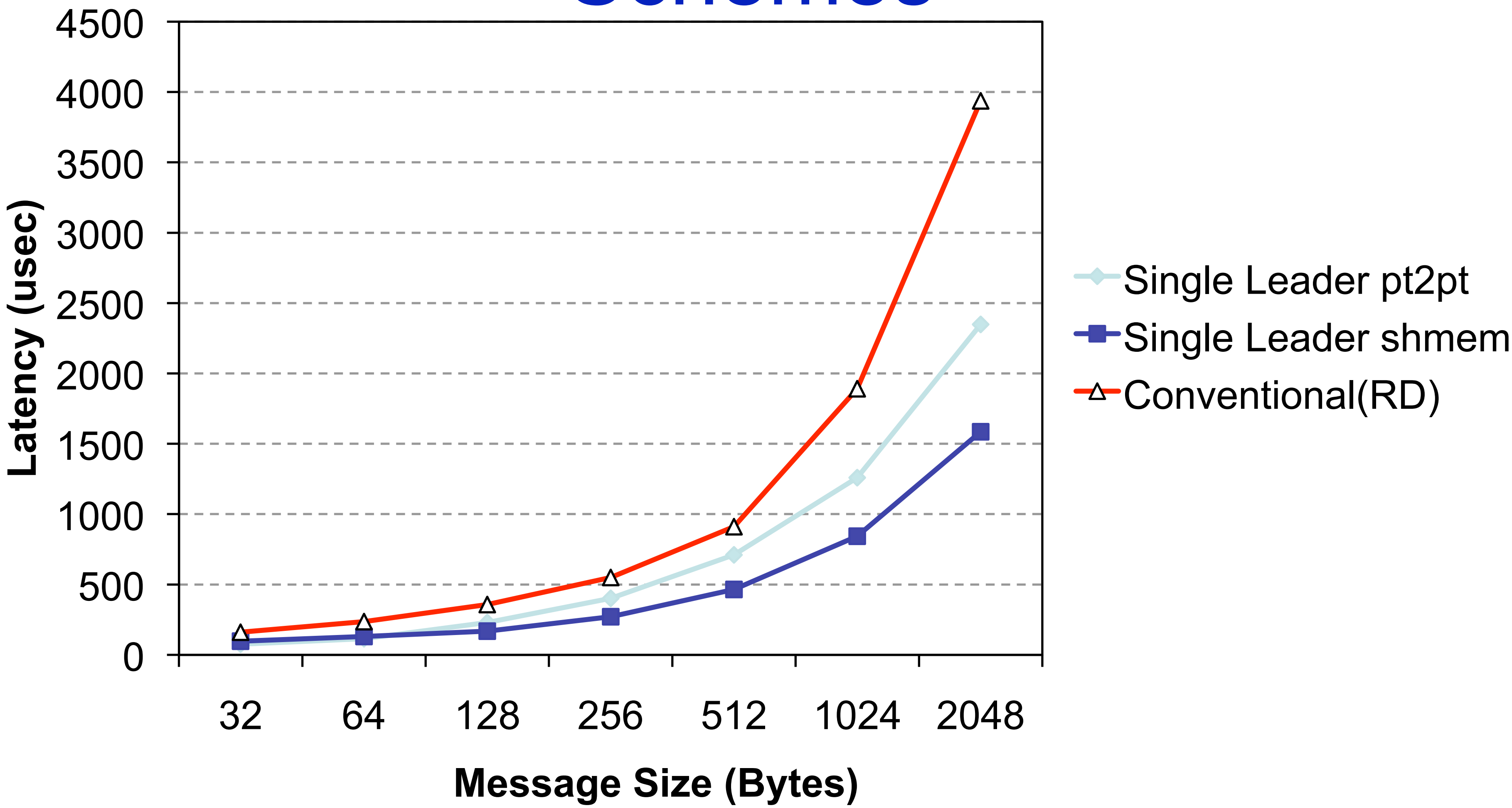


Single Leader Algorithms : Step3 intra-node (shmem)

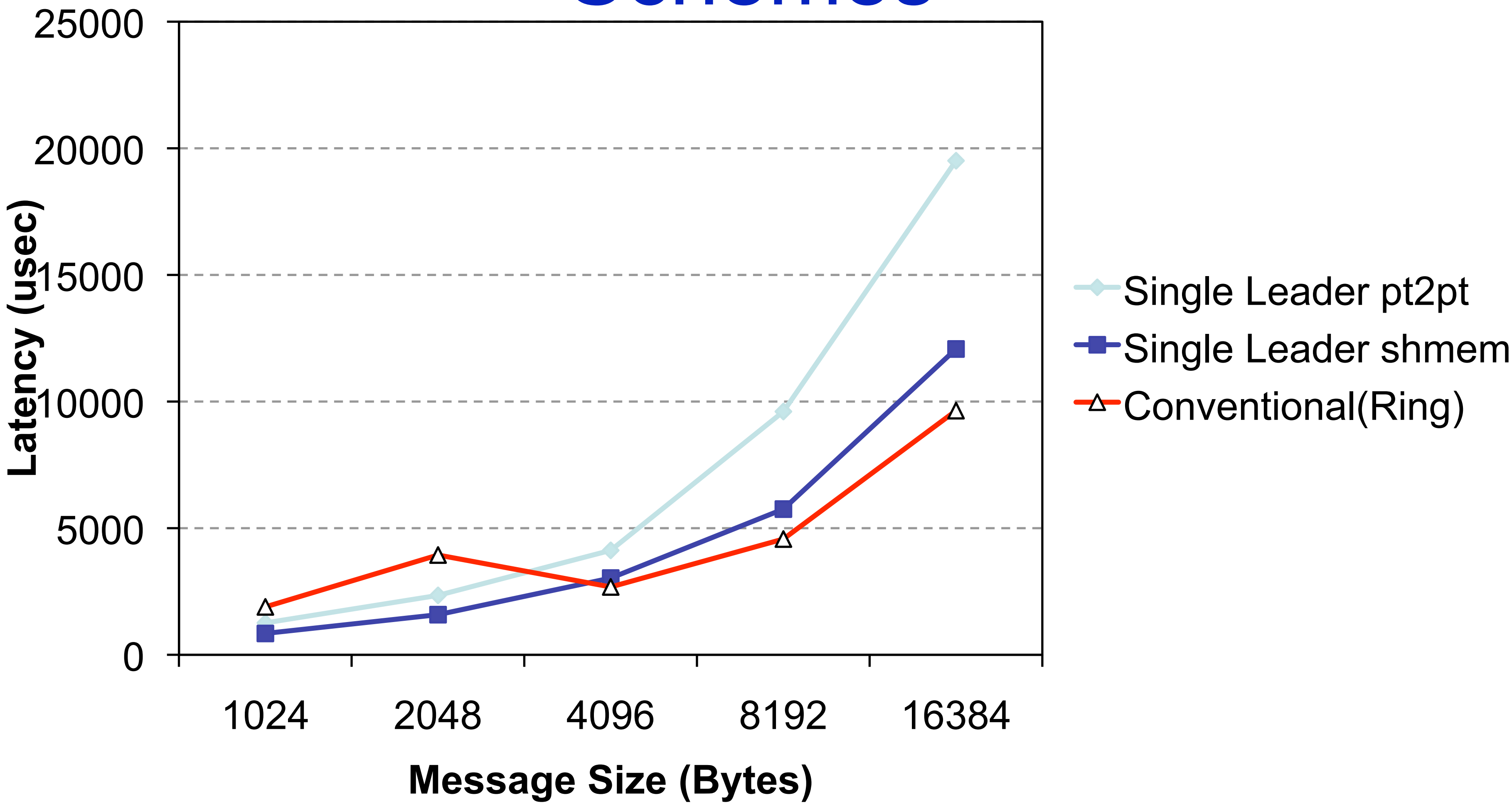


Performance of Single Leader Schemes



Single Leader schemes show potential for improvement

Performance of Single Leader Schemes



Conventional Ring Algorithm performs better for larger messages

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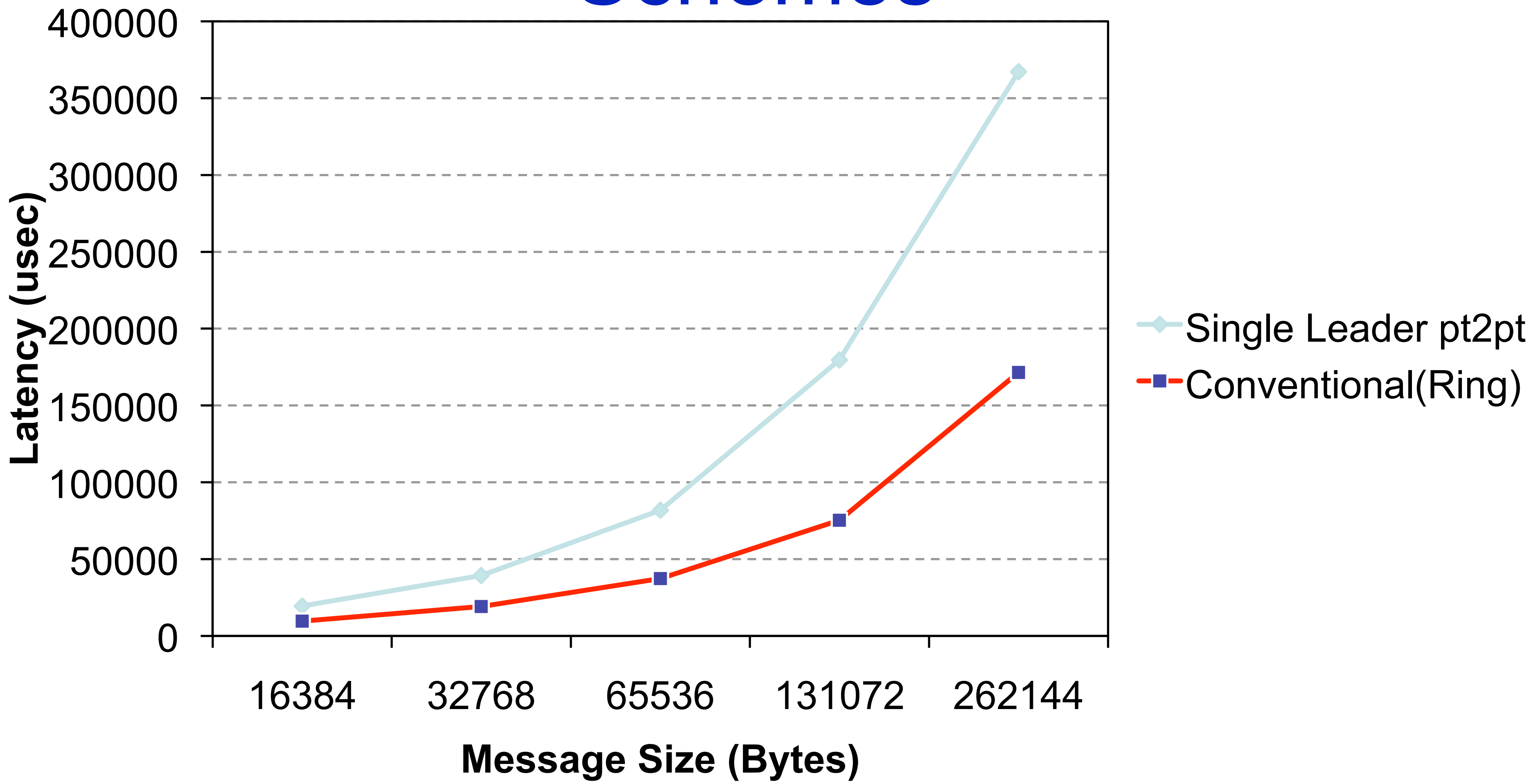

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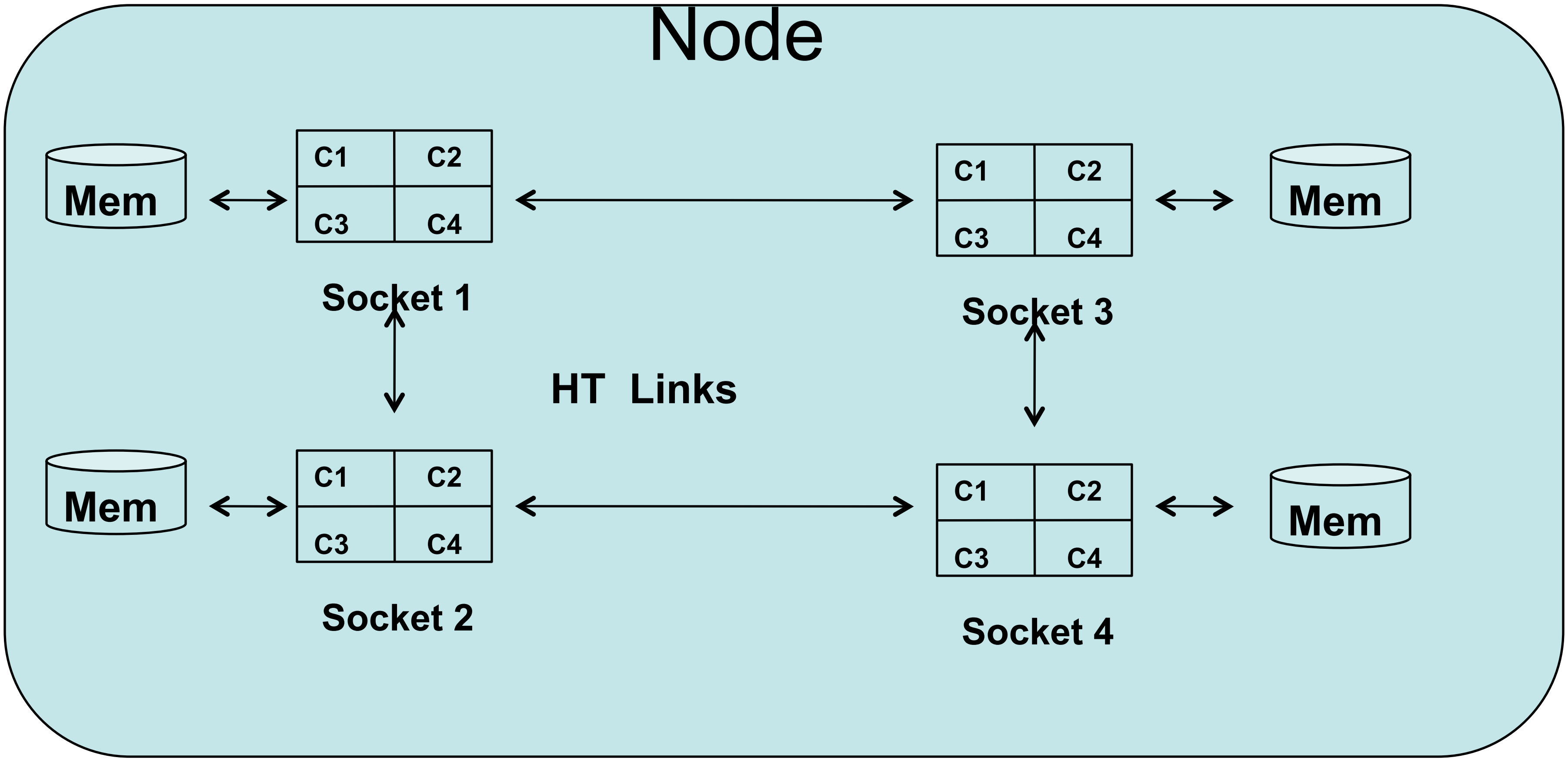
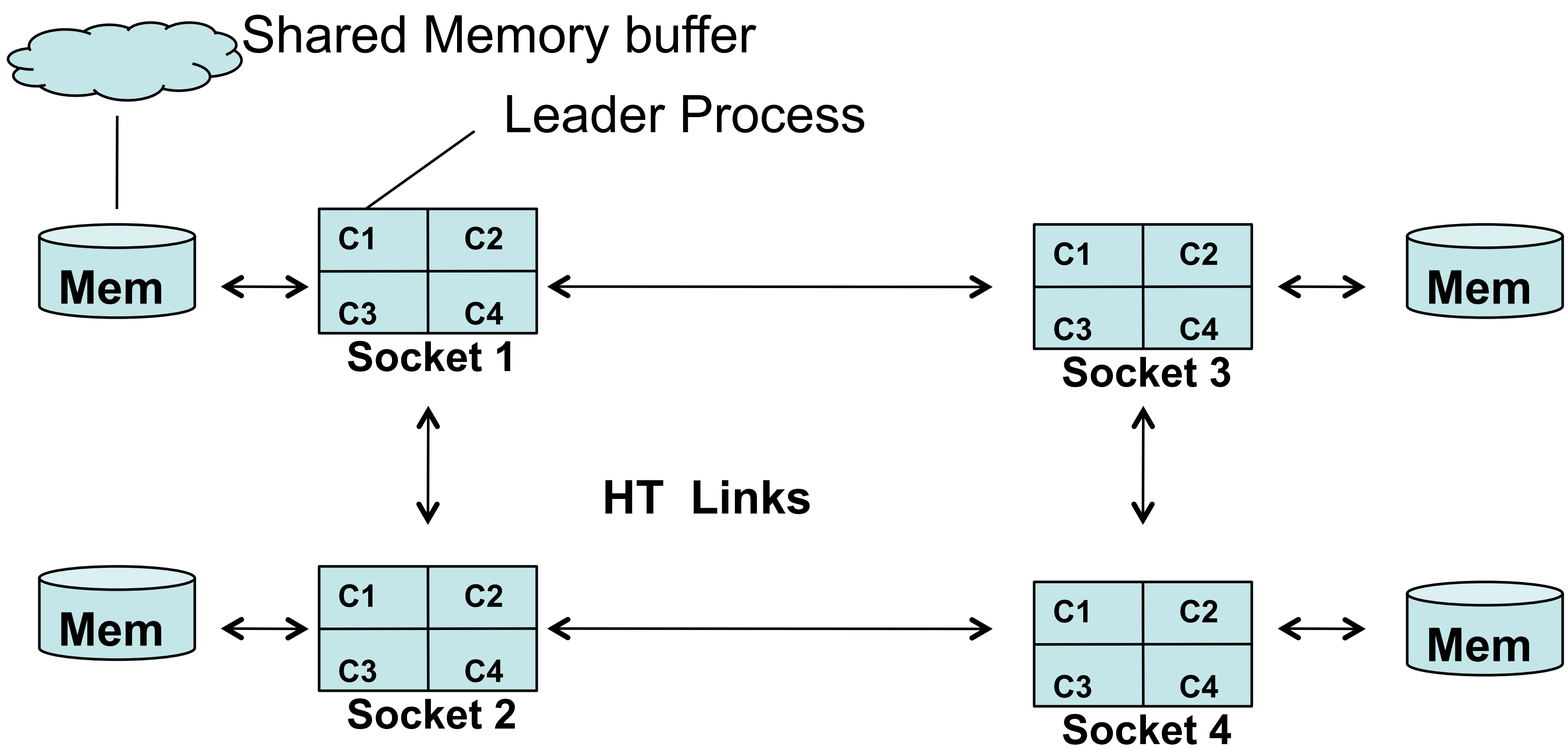
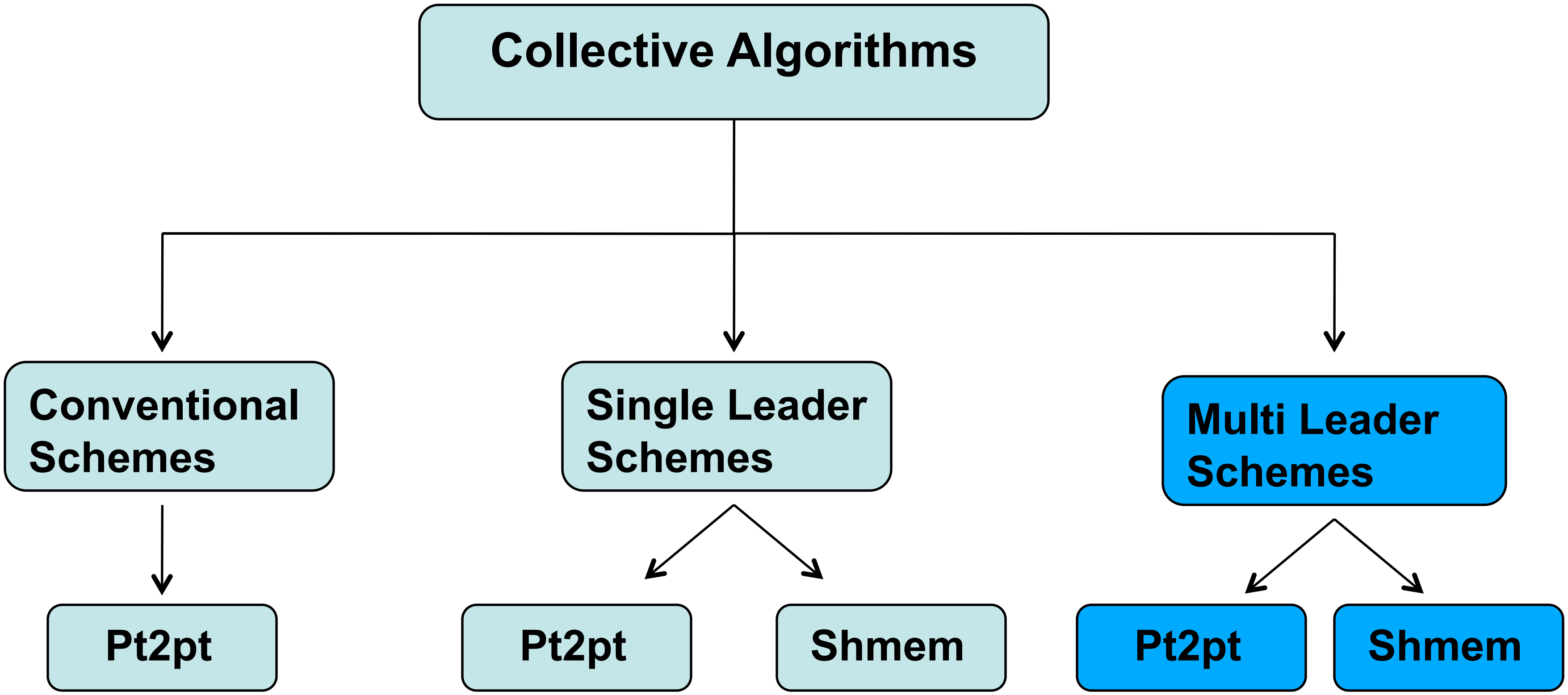
AMD Barcelona Architecture



Single Leader algorithms on the AMD Barcelona Architecture



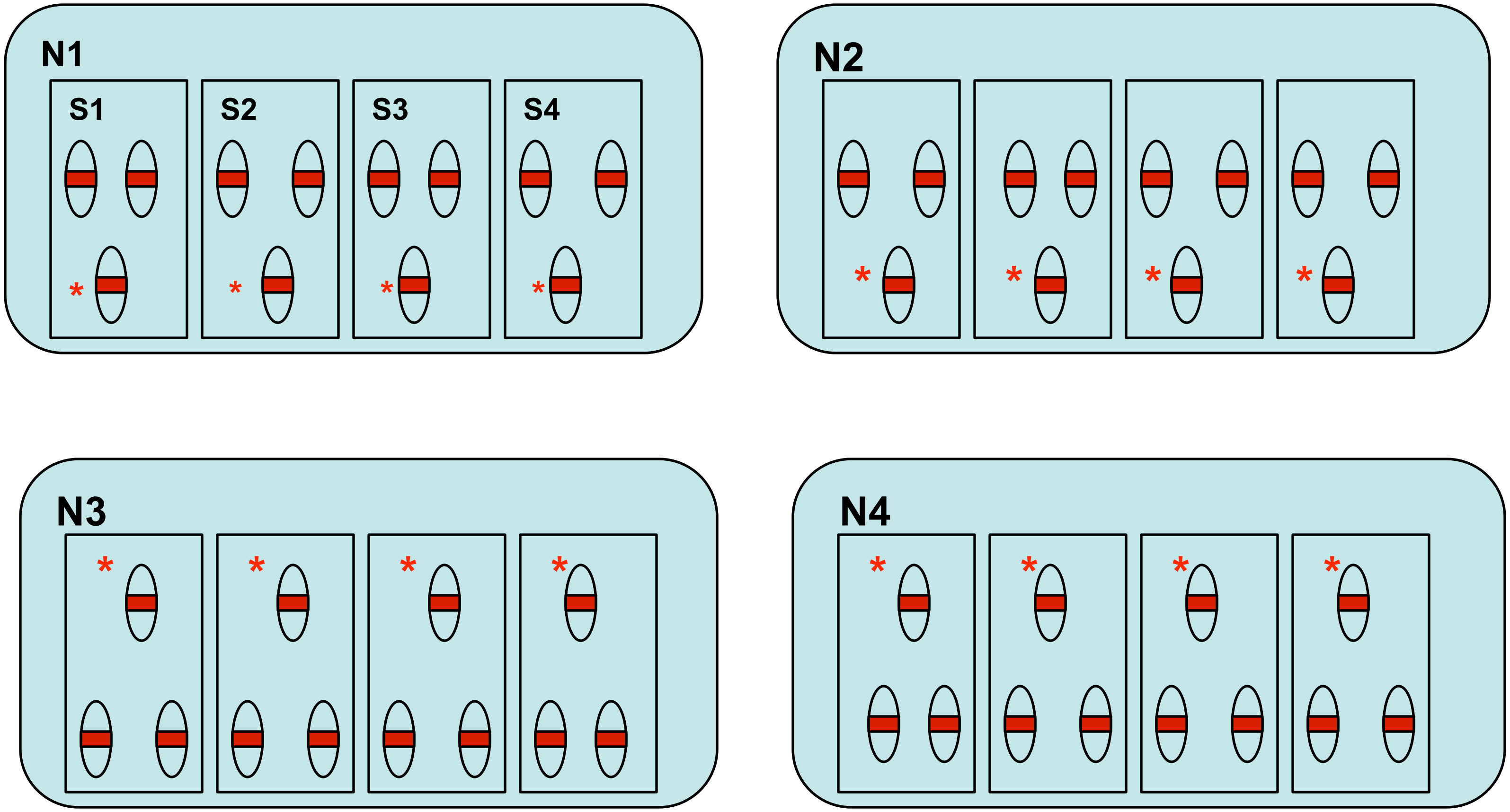
Proposed Collective Design Framework



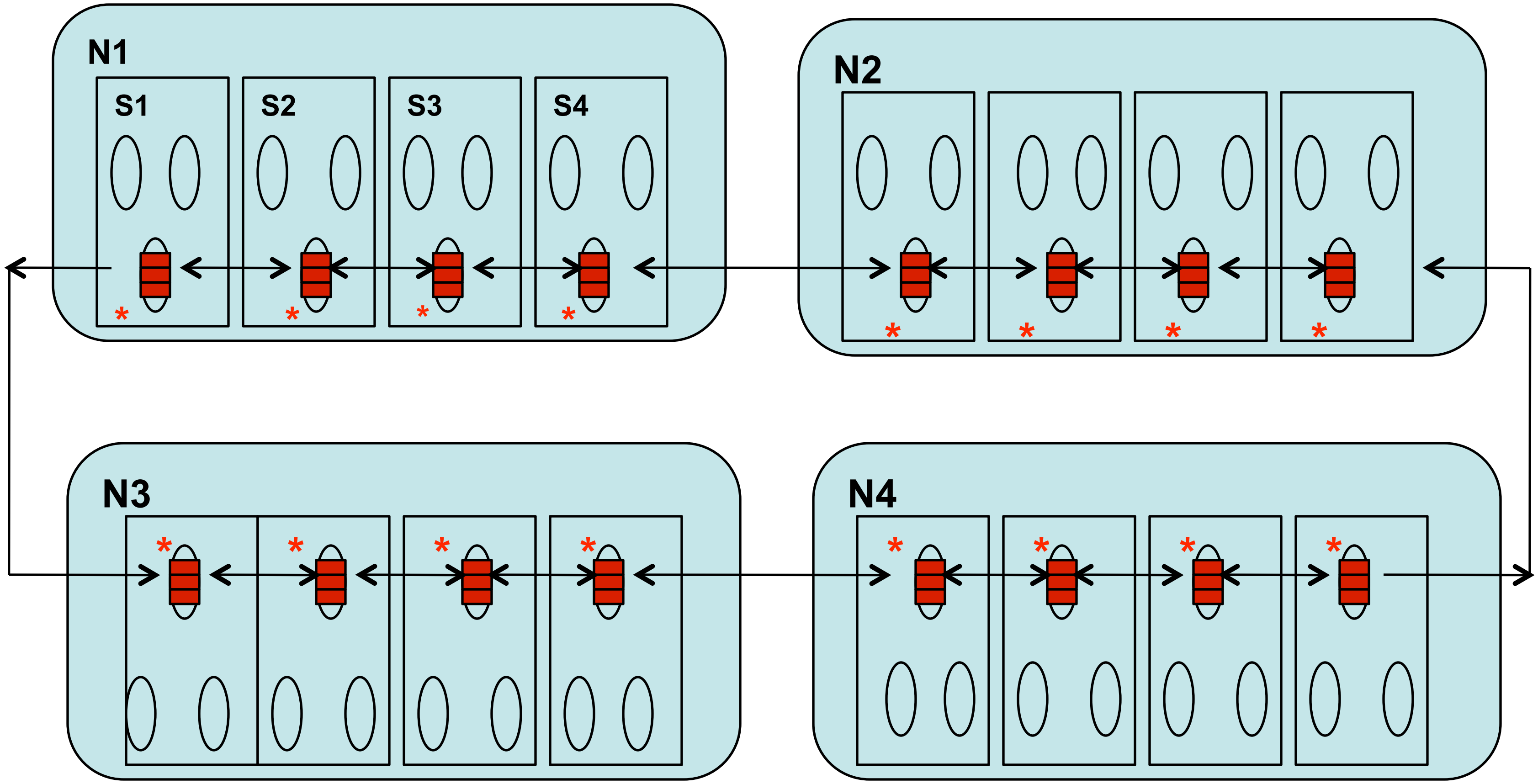
Multi-Leader based Algorithms

- Number of leader processes per node
- Intra-socket and Inter-leader exchange algorithms.

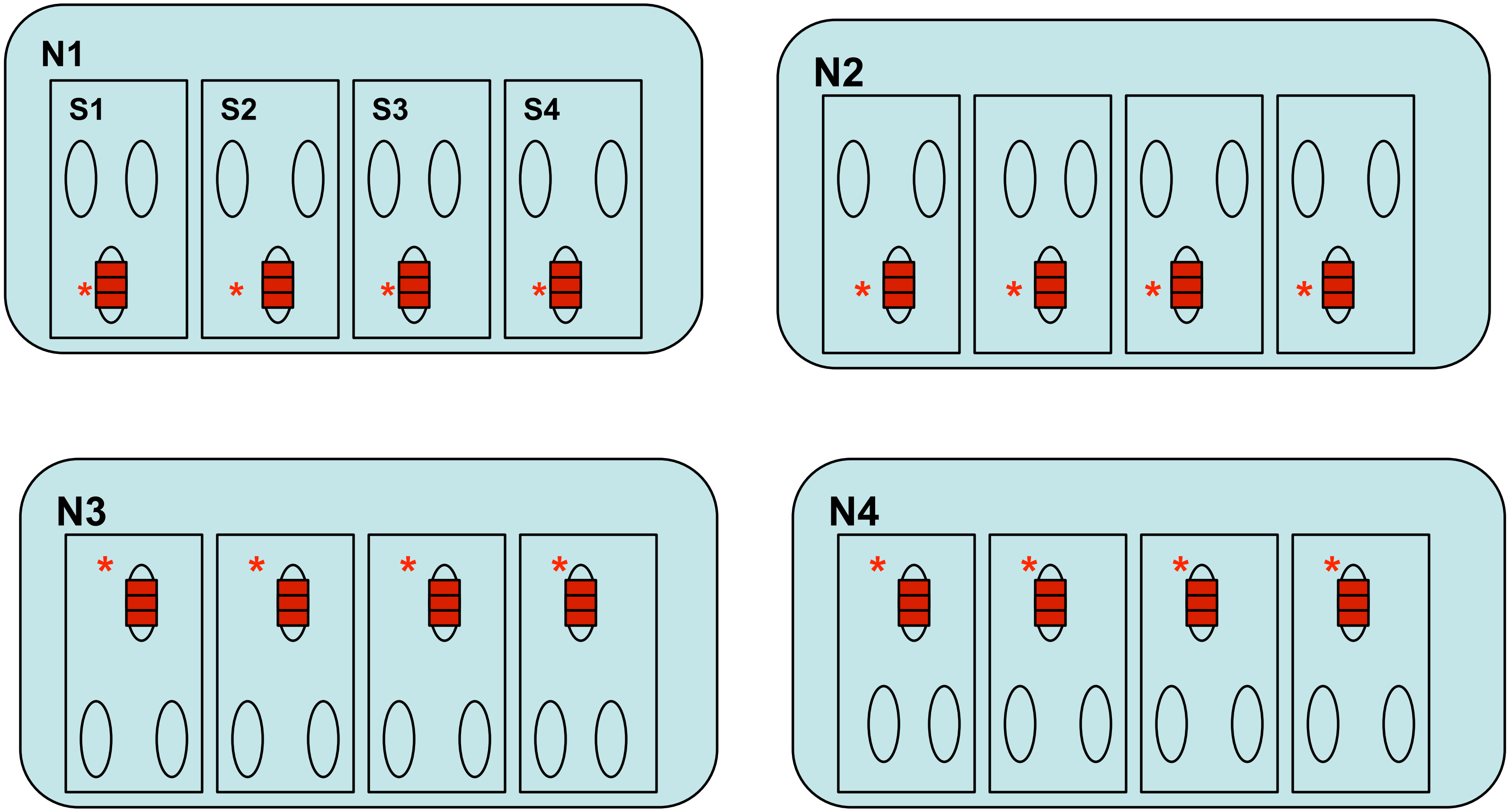
Multi-Leader based Algorithms(Step 1)



Multi-Leader based Algorithms(Step 2)



Multi-Leader based Algorithms(Step 3)



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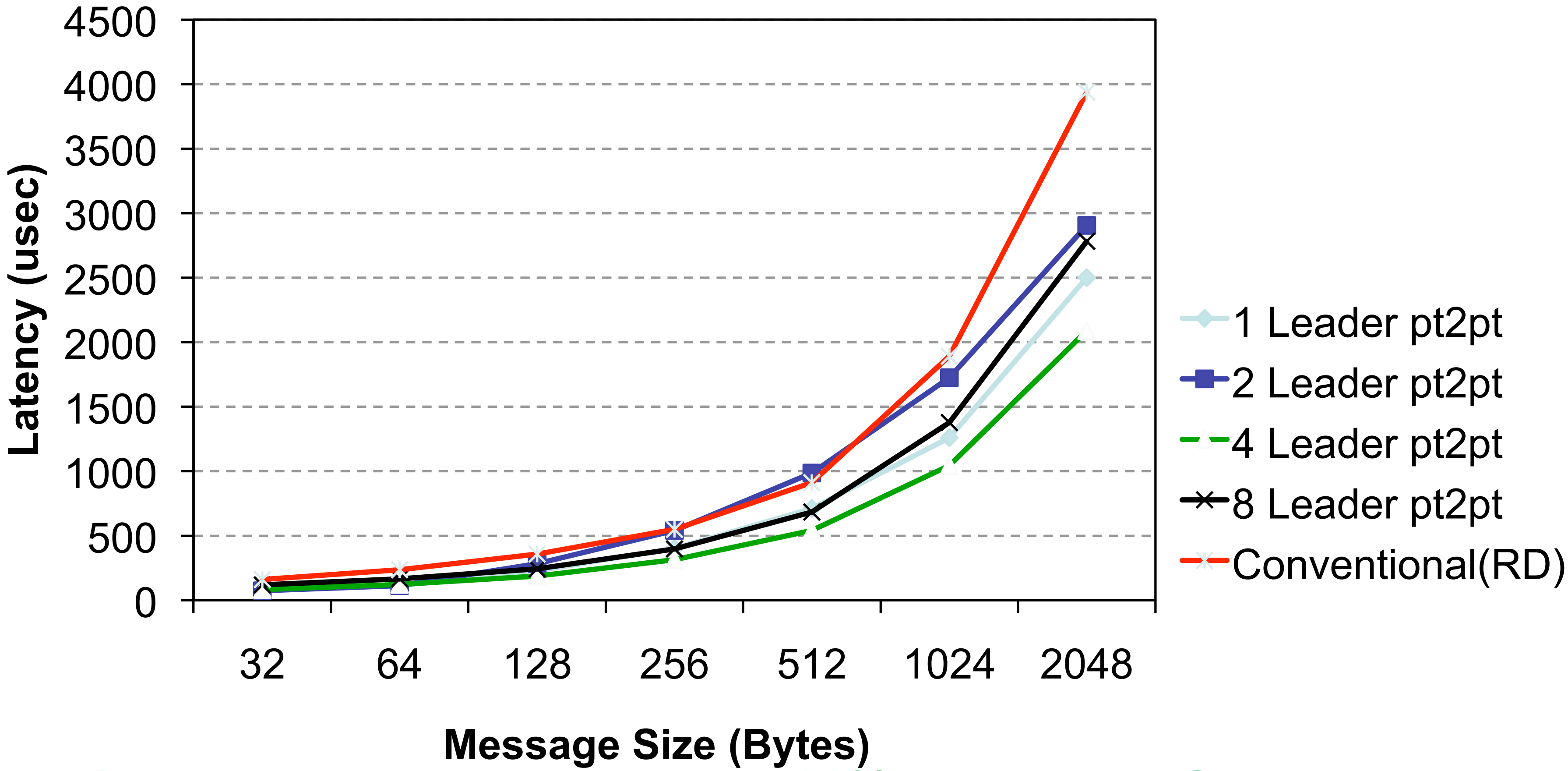
Experimental Test-bed

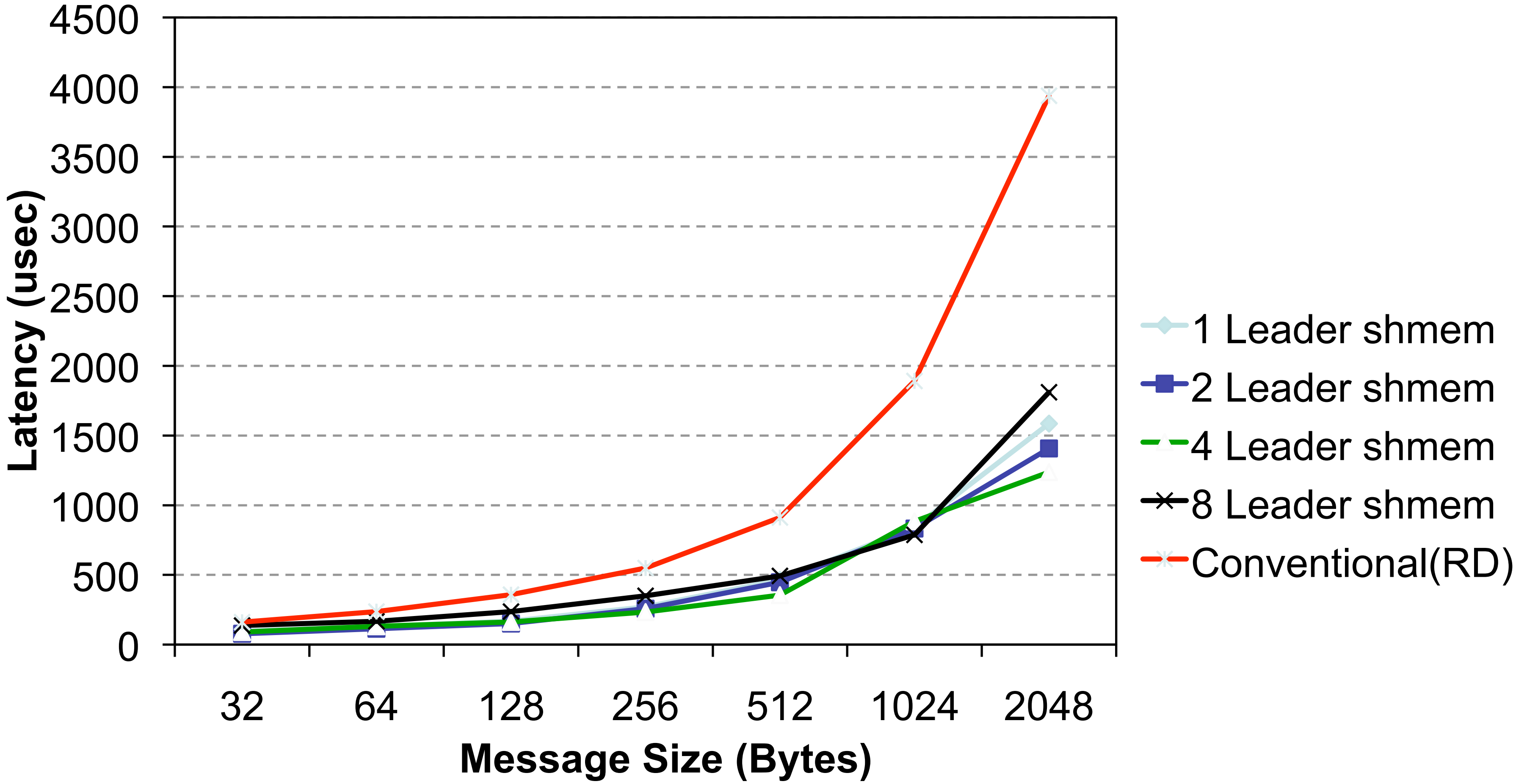
- Each node of our testbed has 16 AMD Opteron 1.95 Ghz processors with 512 KB L2 cache. We used 8 nodes.
- Each node has 16 GB memory and PCI-Express bus, 2 MT25418 DDR HCAs with PCI-Ex interfaces.
- 24-port Mellanox switch is used to connect all the nodes.
- RedHat Enterprise Linux Server 5.

Performance of Multi-Leader pt2pt



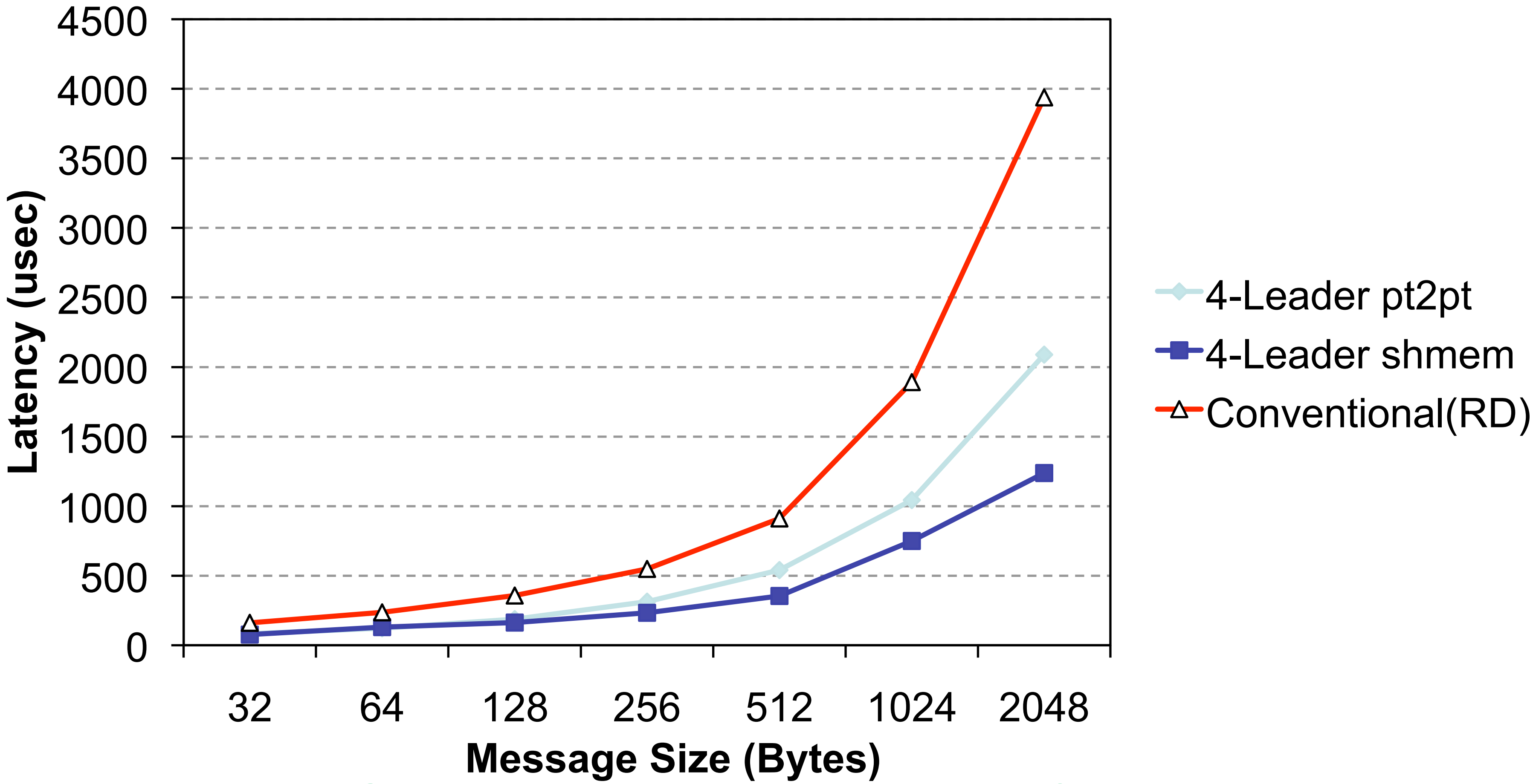
4-Leader scheme does about 20% better than Single Leader scheme and 50% better than RD

Performance of Multi-leader : Shared Memory



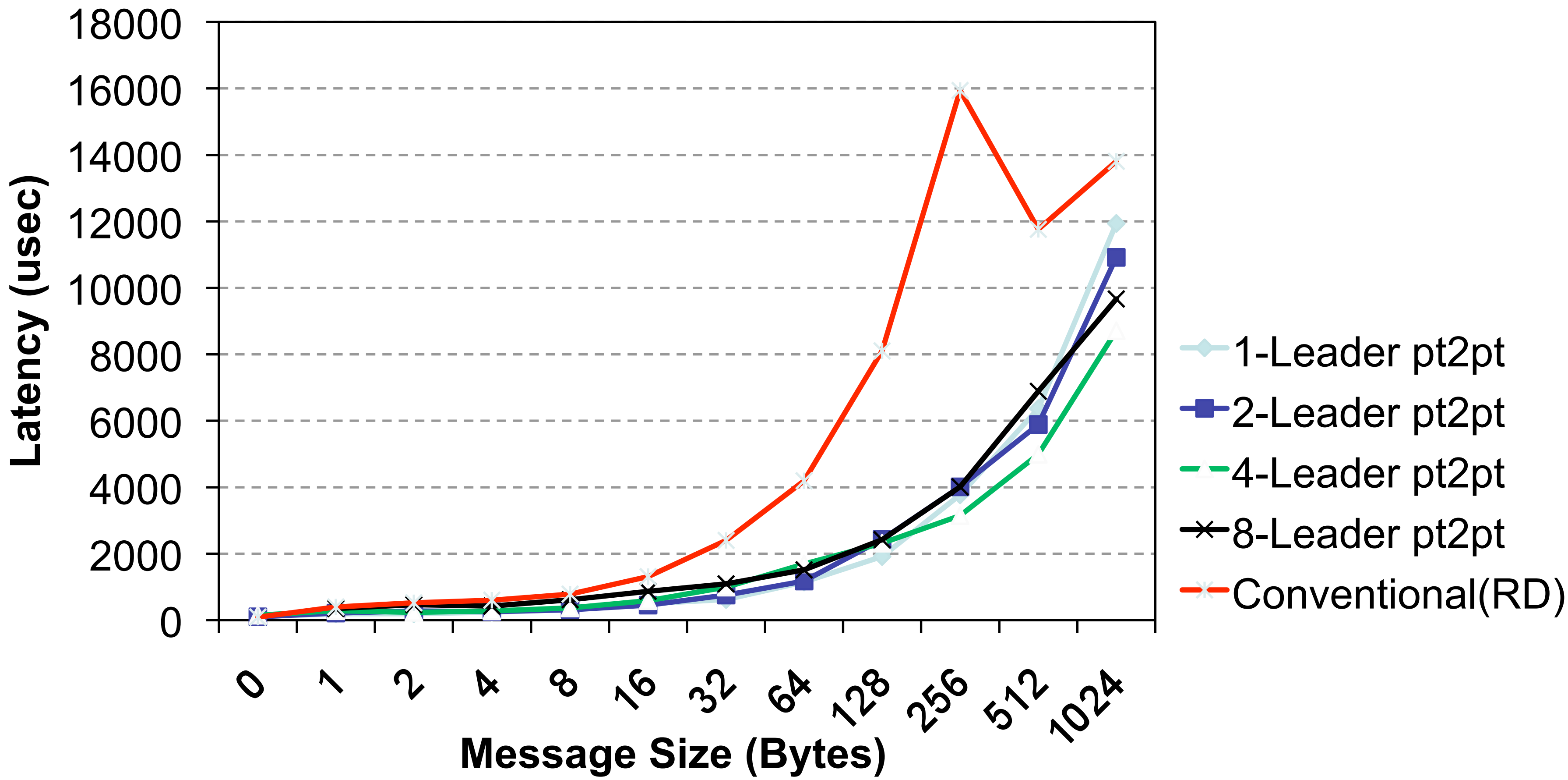
4-Leader scheme performs better than 1-Leader scheme by about 25% and 70% better than RD

Performance of Multi-Leader Schemes (pt2pt Vs Shared Memory)



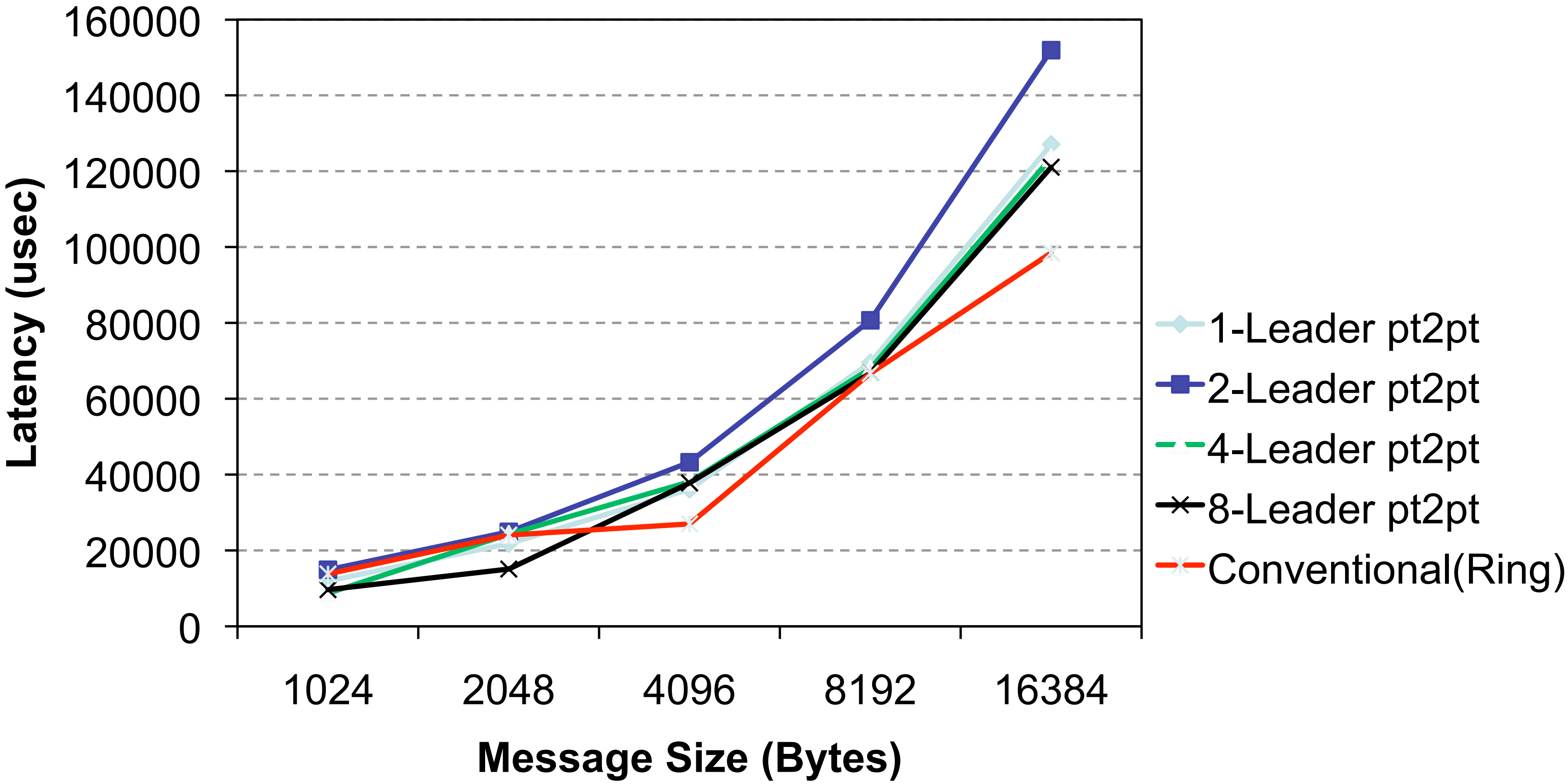
4-Leader Shared Memory approach performs better than 4-Leader Point-to-point scheme by about 40%

Performance of Multi-Leader schemes on large scale Multi-cores



4-Leader Point-to-point scheme outperforms the recursive doubling method on 1024 processes on the TACC Ranger

Performance of Multi-Leader schemes on large scale Multi-cores



Conventional Ring Algorithm performs better for larger messages

Proposed Unified Scheme

	Intra-Node Mechanism	Inter-Leader Algorithm	Design
Small Messages	Point-to-Point	Recursive Doubling	Hierarchical
Medium Messages	Shared Memory	Recursive / Ring	Hierarchical
Large Messages	Point-to-Point	Ring	Conventional

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Conclusions & Future Work

- Single Leader schemes are limited by scalability and memory contention. Proposed Multi-Leader schemes perform show significant performance benefits.
- *Future work:*

  - Examine the benefits of using kernel based zero-copy intra-node exchanges for large messages.

  - A frame-work that can choose leaders in an optimal manner for emerging multi-core systems.

  - Evaluate the impact of such designs on real-world applications.

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Thank you !

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